

CÓRDOBA U<mark>NIVERSIT</mark>Y

SUPERIOR POLYTECHNIC SCHOOL

DEPARTMENT OF COMPUTER SCIENCE AND NUMERICAL ANALYSIS



ARTIFICIAL INTELLIGENCE LANGUAGES

TECHNICAL ENGINEERING IN MANAGEMENT COMPUTER SCIENCE

TECHNICAL ENGINEERING IN SYSTEMS COMPUTER SCIENCE

SECOND YEAR

FIRST FOUR-MONTH PERIOD

ACADEMIC YEAR: 2009 - 2010



ARTIFICIAL INTELLIGENCE LANGUAGES

First part: Scheme

Second part: Prolog

Subject 1.- Introduction to Scheme language Subject 2.- Expressions and Functions Subject 3.- Conditional Predicates and Sentences Subject 4.- Iteration and Recursion Subject 5.- Compound Data Types Subject 6.- Data Abstraction Subject 7.- Reading and Writing

Subject 8.- Introduction to Prolog language

Subject 9.- Basic Elements of Prolog

Subject 10.- Lists

Subject 11.- Re-evaluation and the "cut"

Subject 12.- Input and Output

PROGRAM

First part: Scheme

Subject 1.- Introduction to Scheme language Subject 2.- Expressions and Functions Subject 3.- Conditional Predicates and Sentences Subject 4.- Iteration and Recursion Subject 5.- Compound Data Types Subject 6.- Data Abstraction Subject 7.- Reading and Writing PROGRAM

Contents

- 1. Fundamental Characteristics of Functional Programming
- 2. Historic Summary of Scheme

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- 1. Fundamental Characteristics of Functional Programming
 - ✓ Functional Programming is a subtype of Declarative Programming

- 1. Fundamental Characteristics of Functional Programming
 - ✓ Declarative Programming (1 / 2)
 - Objective: Problem description

"What" problem must be resolved?

- Notice:
 - It does not mind "how" the problem is resolved
 - It **avoids** the implementation features.



- 1. Fundamental Characteristics of Functional Programming
 - ✓ Declarative Programming (2 / 2)
 - Features
 - Expressivity
 - Extensible: 10% 90% rule
 - Protection
 - Mathematic Elegance
 - > Types:
 - **Functional** or Applicative Programming:
 - Lisp, <mark>Scheme</mark>, Haskell, ...
 - Logic Programming: Prolog

- 1. Fundamental Characteristics of Functional Programming
 - ✓ Principle of the "Pure" Functional Programming

"The **expression value** only <mark>depends on</mark> its **sub-expressions** values,

if such sub-expressions exist ".

✓ *Non* collateral effects

The value of "a + b" only depends on "a" and "b".

- ✓ The function term is used in its mathematical sense.
- ✓ No instructions: programming without assignments
 - The impure Functional programming allows the

"assignment instruction"

- 1. Fundamental Characteristics of Functional Programming
 - ✓ Program structure in Functional Programming
 - The program is a function composed of simpler functions
 - > Function execution:
 - Receives the input data: functions arguments or parameters
 - Evaluates the expressions
 - Returns the Result: computed value of the function



- 1. Fundamental Characteristics of Functional Programming
 - ✓ *Type* of Functional Languages
 - Most of them are interpreted languages
 - Some of them have compiled versions
 - ✓ Memory management
 - Implicit memory management:
 - Memory management is a task of the interpreter.
 - The programmer must not worry about memory management.
 - Garbage collection: task of the interpreter.

In short: the programmer must only worry about the Problem description

Contents

- 1. Fundamental Characteristics of Functional Programming
- 2. Historic Summary of Scheme

- 2. Historic Summary of Scheme
 - ✓ LISP
 - ✓ Compilation versus Interpretation
 - ✓ Lexical (or static) versus dynamical scope
 - ✓ Origin of Scheme

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- 2. Historic Summary of Scheme
 - ✓ LIS₽
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- 2. Historic Summary of Scheme
 - ✓ LIS₽
 - John McCarthy (MIT)
 - "Advice Taker" program:
 - Theoretical basis: Logic Mathematics
 - Objective: Deduction and Inferences
 - LISP: LISt Processing (1956 1958)
 - Second historic language of Artificial Intelligence (after IPL)
 - At present time, second historic language in use (after Fortran)
 - LISP is based on Lambda Calculus (Alonzo Church)
 - Scheme is a dialect of LISP

- 2. Historic Summary of Scheme
 - ✓ LIS₽
 - Functional Programming Characteristics
 - Recursion
 - *Lists*
 - Implicit memory management
 - Interactive and *interpreted* programs
 - Symbolic Programming
 - Dynamically scoped for non local variables

- 2. Historic Summary of Scheme
 - ✓ LIS₽
 - LISP's contributions:
 - Built in functions
 - Garbage collection
 - Definition Formal Language: LISP itself



- 2. Historic Summary of Scheme
 - ✓ LIS₽
 - Applications: Artificial Intelligence Programs
 - Theorem verification and testing
 - Symbolic differentiation and integration
 - Search Problems
 - Natural Language Processing
 - Computer Vision
 - Robotics
 - Knowledge Representation Systems
 - Expert Systems
 - And so on

- 2. Historic Summary of Scheme
 - ✓ LIS₽
 - Dialects (1 /2)
 - Mac LISP (Man and computer or Machine aided cognition): East Coast Version
 - Inter LISP (Interactive LISP): West Coast Version
 - Bolt, Beranek y Newman Company (BBN)
 - Research Center of Xerox at Palo Alto (Texas)
 - LISP Machine



- 2. Historic Summary of Scheme
 - ✓ LIS₽
 - Dialects (2 / 2)
 - Mac LISP (Man and computer or Machine aided cognition): East Coast Version
 - C-LISP: Massachusetts University
 - Franz LISP: California University (Berkeley). Compiled version.
 - NIL (New implementation of LISP): MIT.
 - PSL (Portable Standard LISP): Utah University
 - Scheme: MIT.
 - T(True): Yale University.
 - Common LISP

- 2. Historic Summary of Scheme
 - ✓ LISP
 - ✓ Compilation versus Interpretation
 - ✓ Lexical (or static) versus dynamical scope
 - ✓ Origin of Scheme

- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - > Compilation:
 - The source code (high level) is transformed into executable code (low level), which can be independently run.

- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - > Compilation

Source code \rightarrow



- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - > Compilation





Compilation errors

Subject 1.- Introduction to Scheme language

- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - > Compilation

Source code \rightarrow



→ Executable code

Subject 1.- Introduction to Scheme language

- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - > Compilation

Input data





→ Executable code

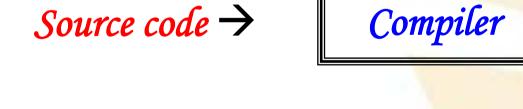
- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - > Compilation

Input data



- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - > Compilation

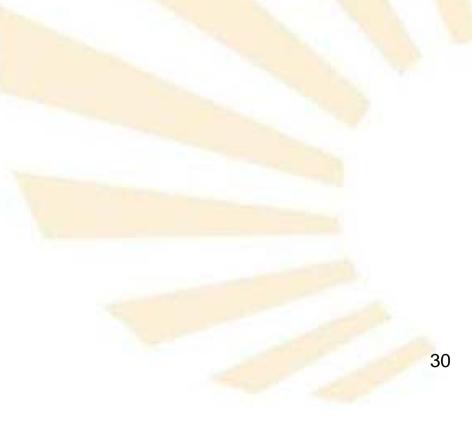
Input data





- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - > Interpretation

- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - Interpretation or simulation: consists of a cycle of three stages



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 - 1. *Analysis*: the source code is analysed to determine the following correct sentence to be run.

- 2. Historic Summary of Scheme
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 - 1. *Analysis*: the source code is analysed to determine the following correct sentence to be run.
 - 2. **Generation**: the sentence is transformed into executable code.

- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - Interpretation or simulation: consists of a cycle of three stages
 - 1. Analysis: the source code is analysed to determine the following correct sentence to be run.
 - 2. *Generation*: the sentence is transformed into executable code.
 - 3. Execution: the executable code is run.

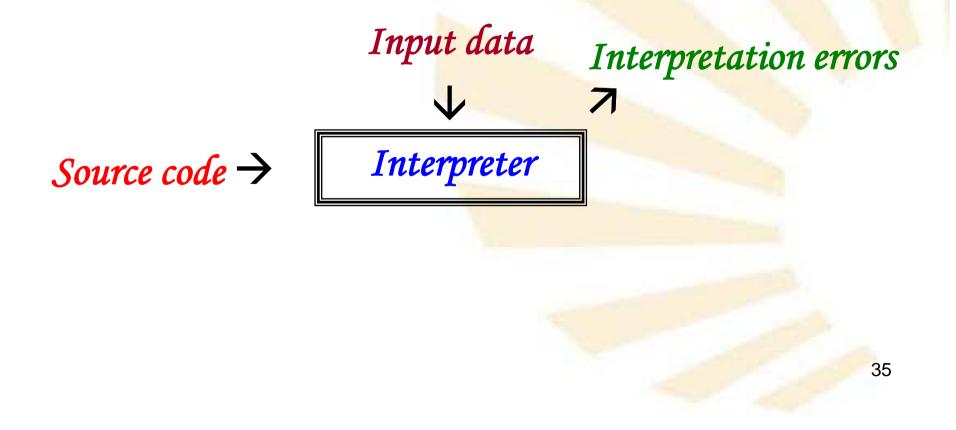
- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation

Interpreter

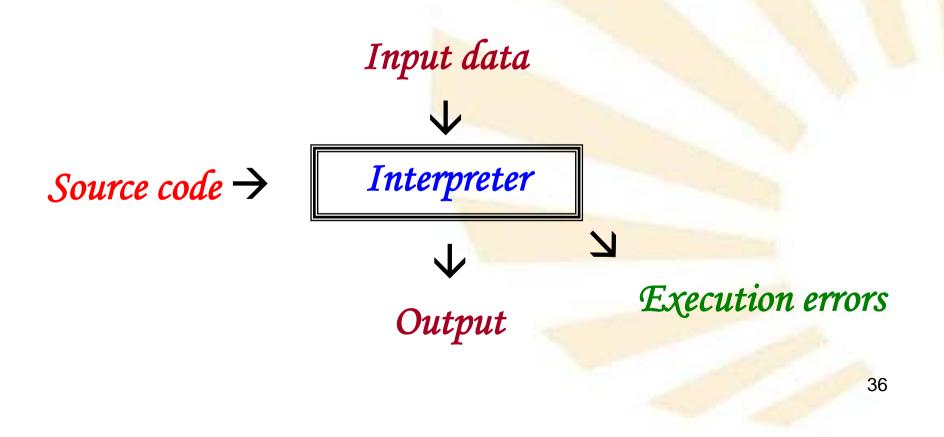
> Interpretation



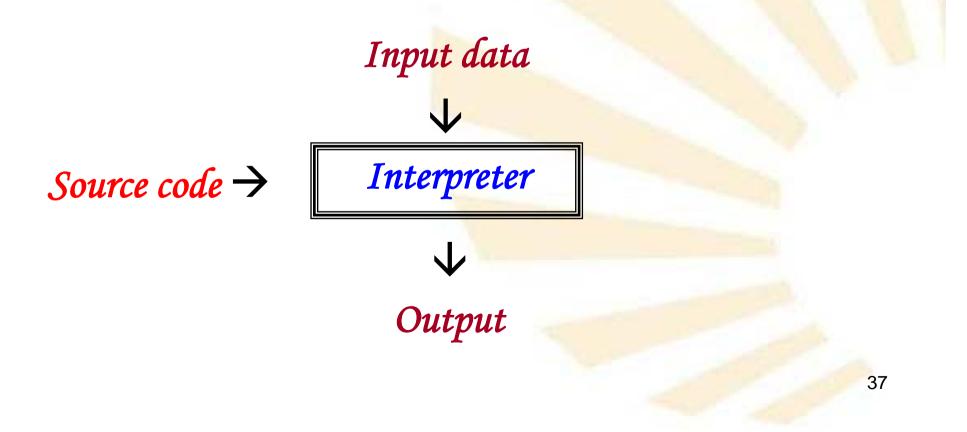
- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
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- 2. Historic Summary of Scheme
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- 2. Historic Summary of Scheme
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- 2. Historic Summary of Scheme
 - ✓ Compilation versus interpretation
 - Compilation
 - Independent
 - Memory necessities
 - Efficient
 - Global
 - No interaction
 - Closed code during execution

- Interpretation
 - <mark>Depen</mark>dent
 - No memory necessities
 - Less efficient
 - Local
 - Interaction
 - Open code during execution

- 2. Historic Summary of Scheme
 - ✓ LISP
 - ✓ Compilation versus Interpretation
 - ✓ Lexical (or static) versus dynamical scope
 - ✓ Origin of Scheme

- 2. Historic Summary of Scheme
 - ✓ Lexical (or static) versus dynamical scope
 - The scope rules determine the declaration of non local identifiers
 - > Non local identifiers:
 - Variables or functions which can be used in a function or procedure but are not declared in that function or procedure
 - Two types
 - Lexical or static scope
 - With "blocks structure": Pascal, Scheme
 - Without "blocks structure": C, Fortran
 - Dynamical scope:
 - Always with "blocks structure": Lisp, SNOBOL, APL

- 2. Historic Summary of Scheme
 - ✓ Lexical (or static) versus dynamical scope
 - Block structure
 - A procedure or function can call
 - Itself
 - Its children (but not its grandchildren...)
 - Its brothers (but not its nephews)
 - Its father, grandfather, great-grandfather, ...
 - The brothers of its father, grandfather, ...
 - A procedure or function can be called by
 - Itself
 - Its father (but not by its grandfather, ...)
 - Its children, grandchildren, great-grandchildren, ...
 - Its brothers and their children, grandchildren, ...

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Example of blocks structure

Declaration of procedure f

Declaration of procedure g

Declaration of procedure h

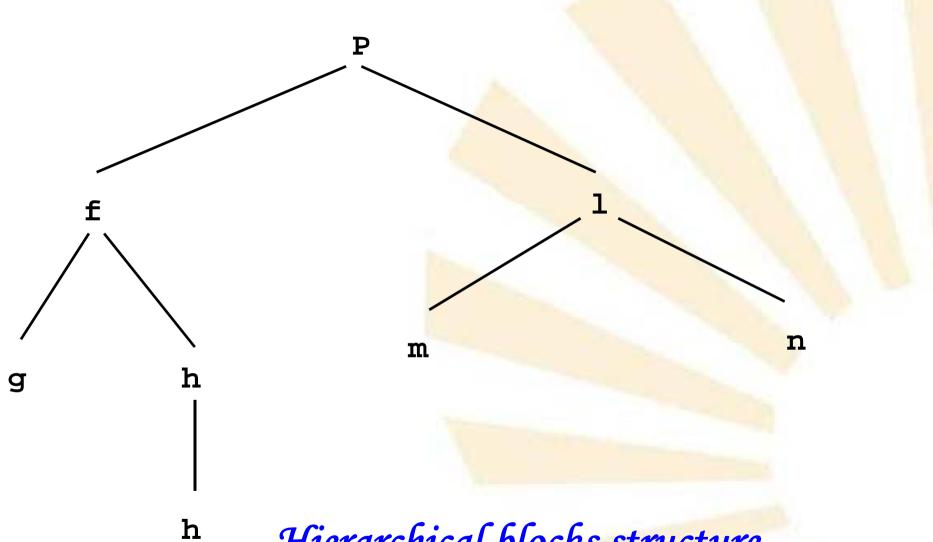
Declaration of procedure k

Declaration of procedure 1

Declaration of procedure m

Declaration of procedure n

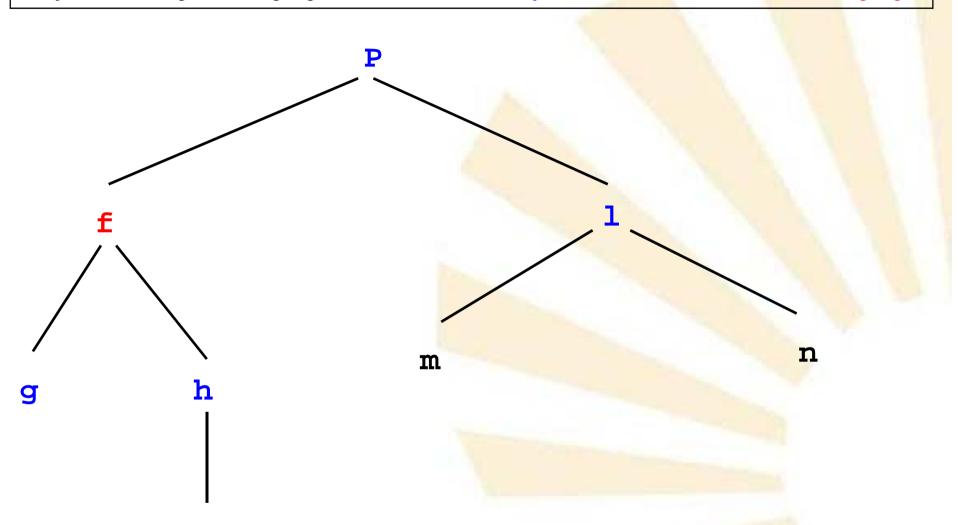
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Hierarchical blocks structure

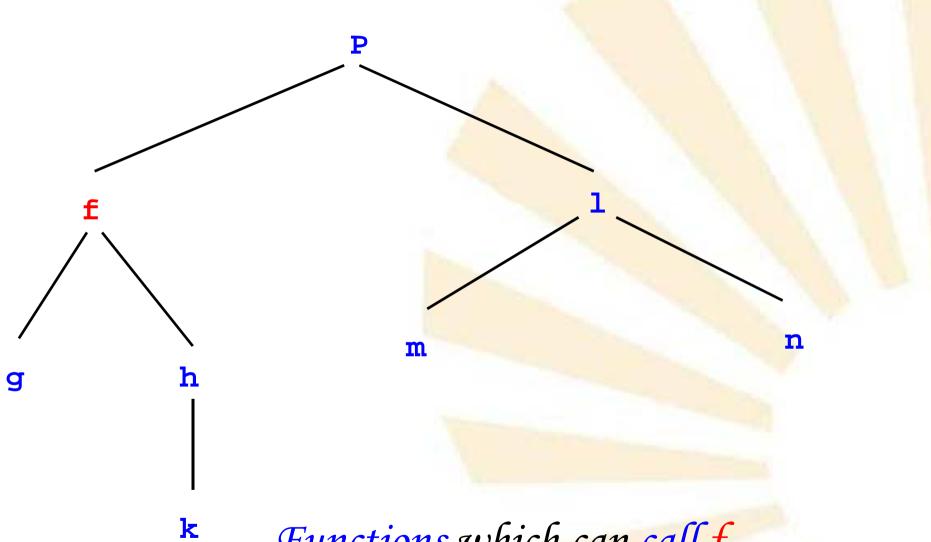
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Functions which can be called by f

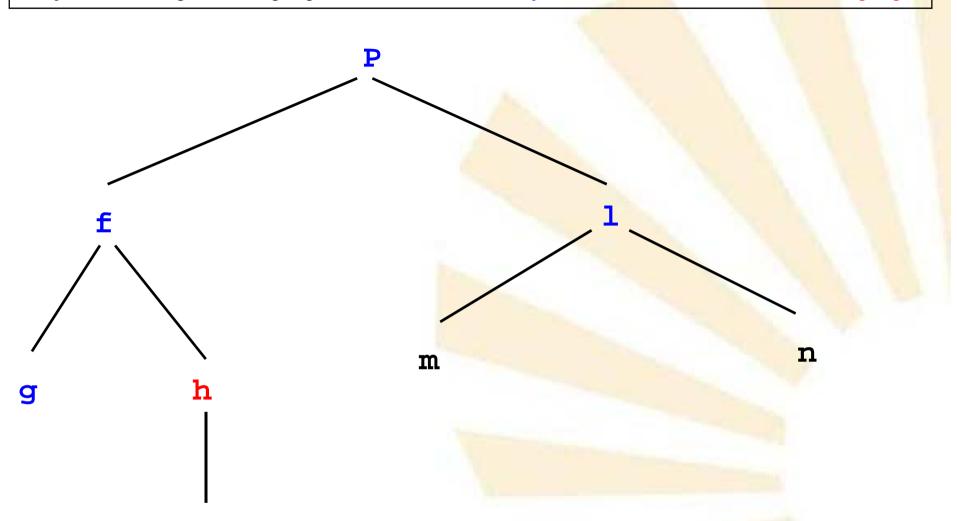
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Functions which can call f

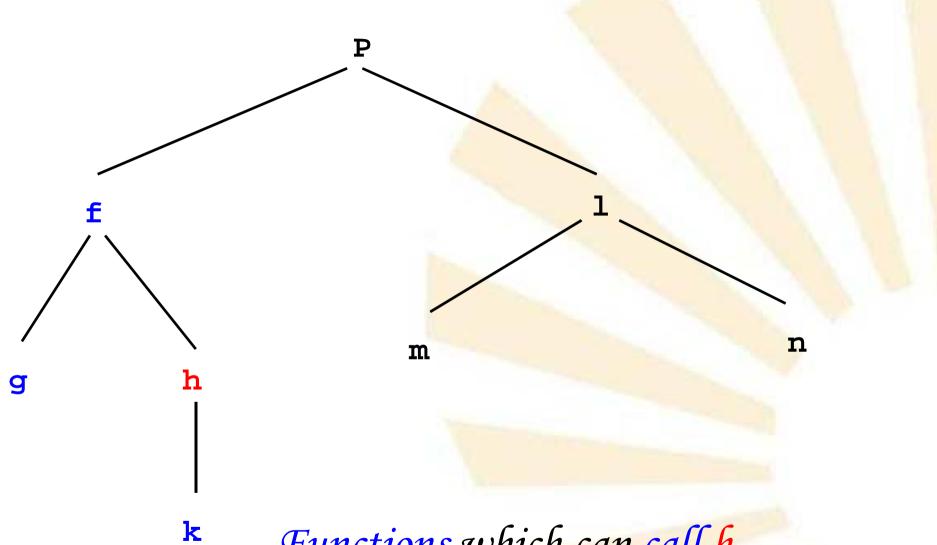
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Functions which can be called by h

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Functions which can call h

- 2. Historic Summary of Scheme
 - ✓ Lexical (or static) versus dynamical scope
 - Lexical or static scope
 - The declaration of a non local identifier depends on the closest lexical context:

You only have to **read** the program

to determine the declaration of an identifier.

- The closest nesting rules:
 - The **scope** of a procedure (*) **f** includes the procedure **f**.
 - If a non local identifier x is used in f then the declaration of x must be found in the closest procedure g which includes f
 - Notice (*): procedure, function or block

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Example: Lexical scope

Lexical scope with "block structure"

Declaration of procedure h Declaration of variable x Declaration of variable y Declaration of variable z	(y1)
Declaration of procedure g Declaration of variable x Declaration of variable y	(y2)
Declaration of procedure f Declaration of variable	
Use of x $(\rightarrow x3)$ Use of y $(\rightarrow y2)$ Use of z $(\rightarrow z1)$	
Use of x $(\rightarrow x2)$ Use of y $(\rightarrow y2)$ Use of z $(\rightarrow z1)$ Call to f	
Use of x (→ x1) Use of y (→ y1) Use of z (→ z1) Call to g	

- 2. Historic Summary of Scheme
 - ✓ Lexical (or static) versus dynamical scope
 - Lexical or static scope
 - Without block structure:
 - If x is not local for a specific function then it is not local for all functions

Example in C: without "block structure"

```
int x; /* x1 */ Global variables are
int y; /* y1 */
                      not recommended
int z; /* z1 */
main()
    int x; /* x2 */
    int y; /* y2 */
    /* Use of x \rightarrow x^2 */
    /* Use of y \rightarrow y2 */
    /* Use of z \rightarrow z1 */
    /* Call to f */
    f ();
 f()
   int x; /* x3 */
    /* Use of x \rightarrow x3 */
    /* Use of y \rightarrow y1 */
    /* Use of z \rightarrow z1 */
                                         51
```

- 2. Historic Summary of Scheme
 - ✓ Lexical (or static) versus dynamical scope
 - > Dynamical scope:
 - The declaration of an identifier depends on the execution of the program

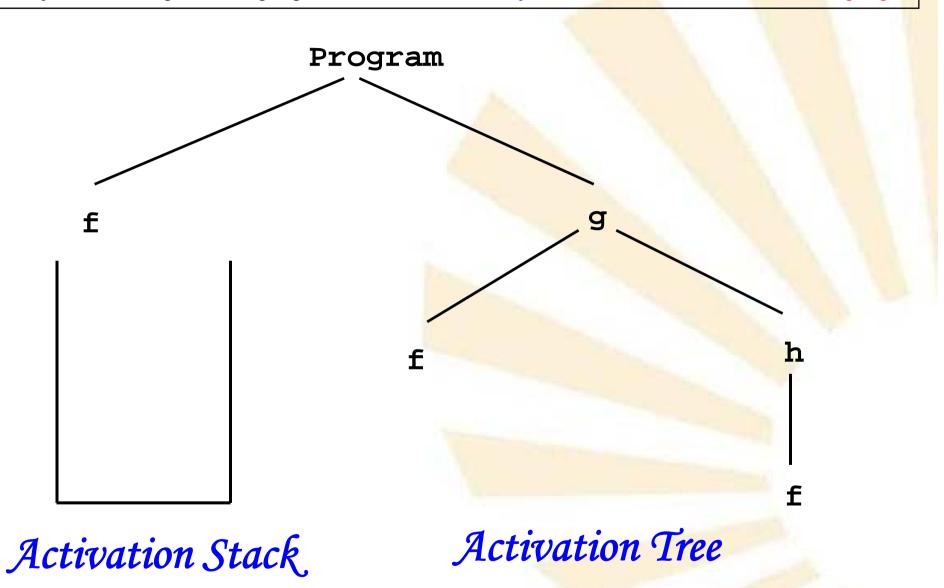
You have to **run** the program to determine the declaration of an identifier

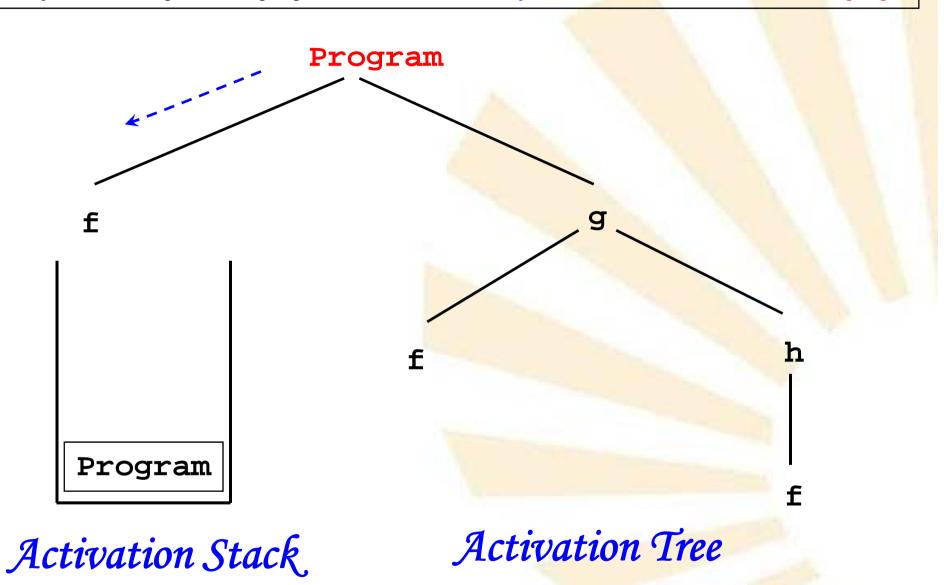
- The closest activation rules:
 - The scope of a procedure (*) f includes the procedure f.
 - If a non local identifier x is used in the activation of f then the declaration of x must be found in the closest active procedure g with a declaration of x
 - Notice (*): procedure, function or block

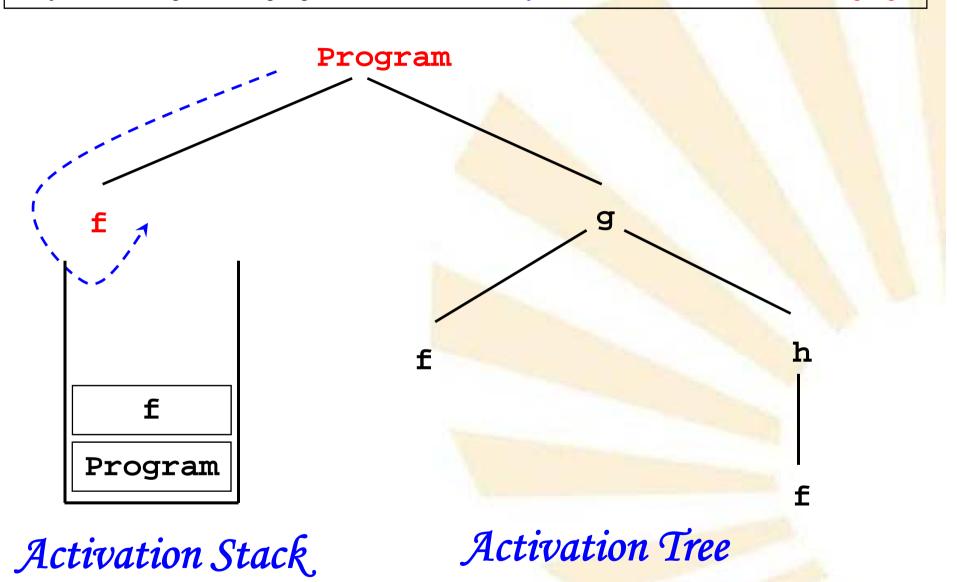
- 2. Historic Summary of Scheme
 - ✓ Lexical (or static) versus dynamical scope
 - > Notice:
 - The dynamical scope allows that an identifier can be associated to different declarations during the program execution

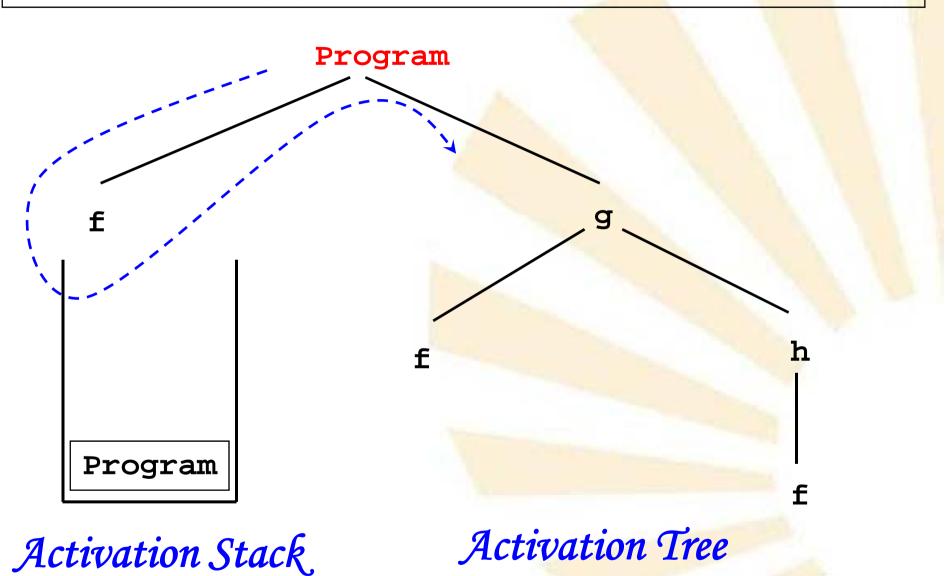
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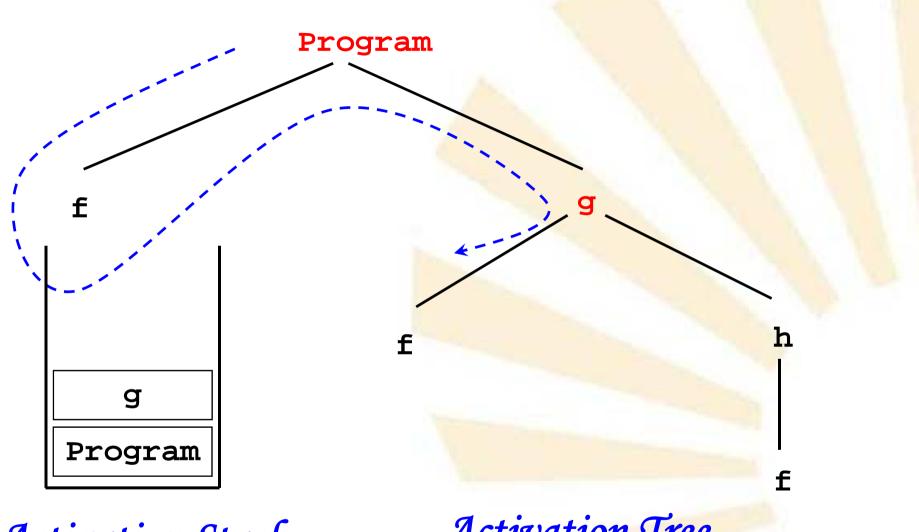
Program **Declaration** of variable **x** Example: Lexical Declaration of procedure f Use of x Declaration of procedure g **Declaration** of variable x versus Declaration of procedure h Dynamical Use of x Call to f scope Call to f Call to h if condition = true then Call to q else Use of x Use of x Call to f Call to g



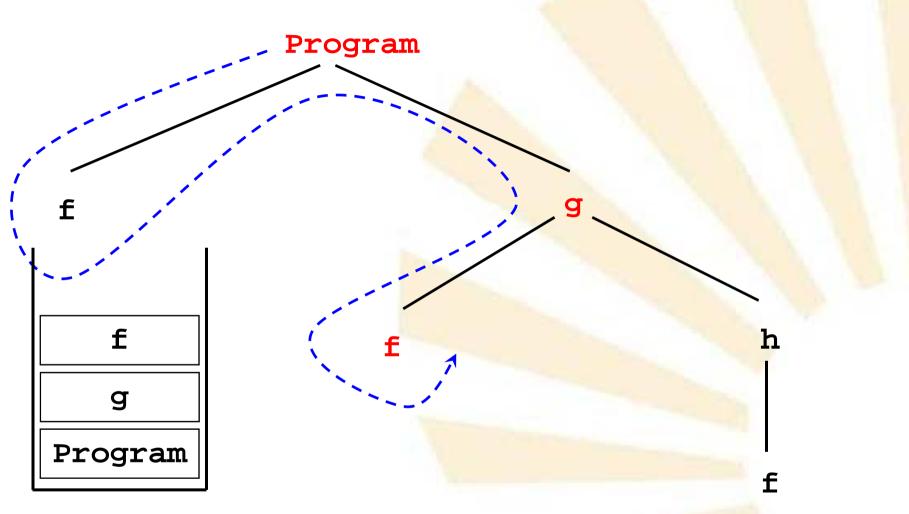




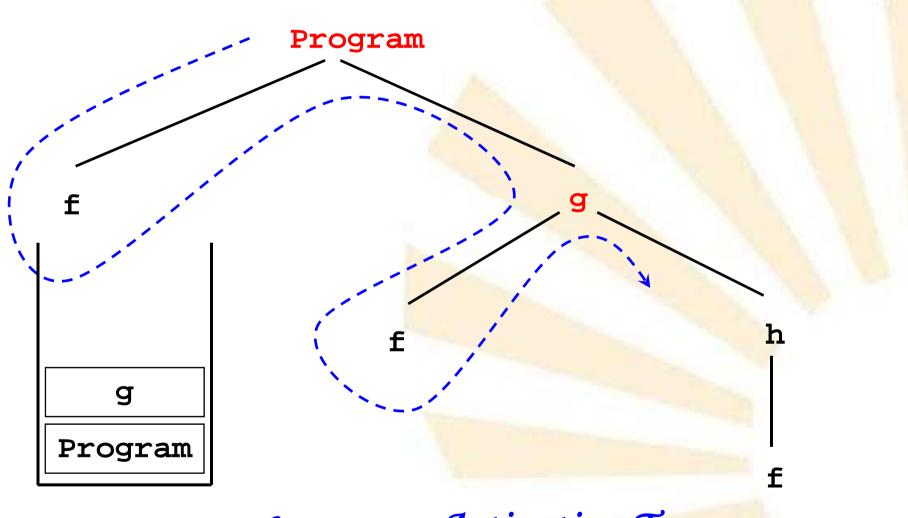




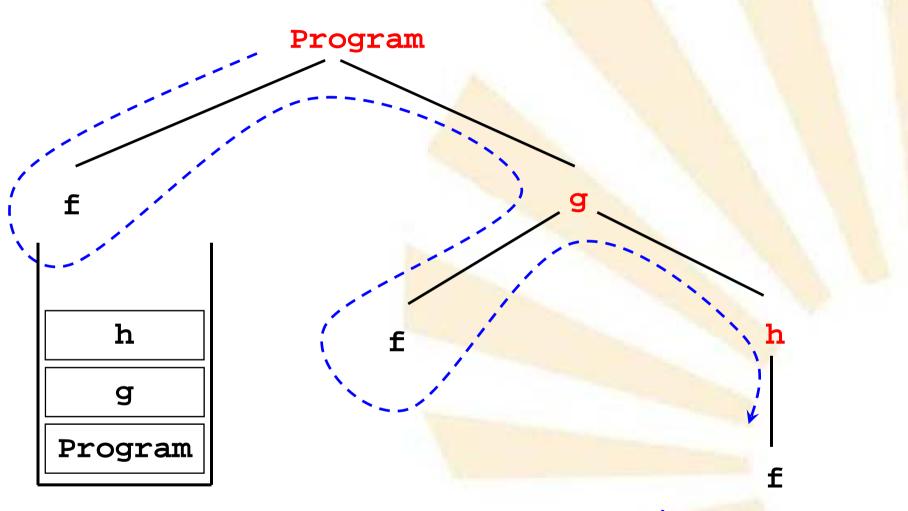
Activation Stack



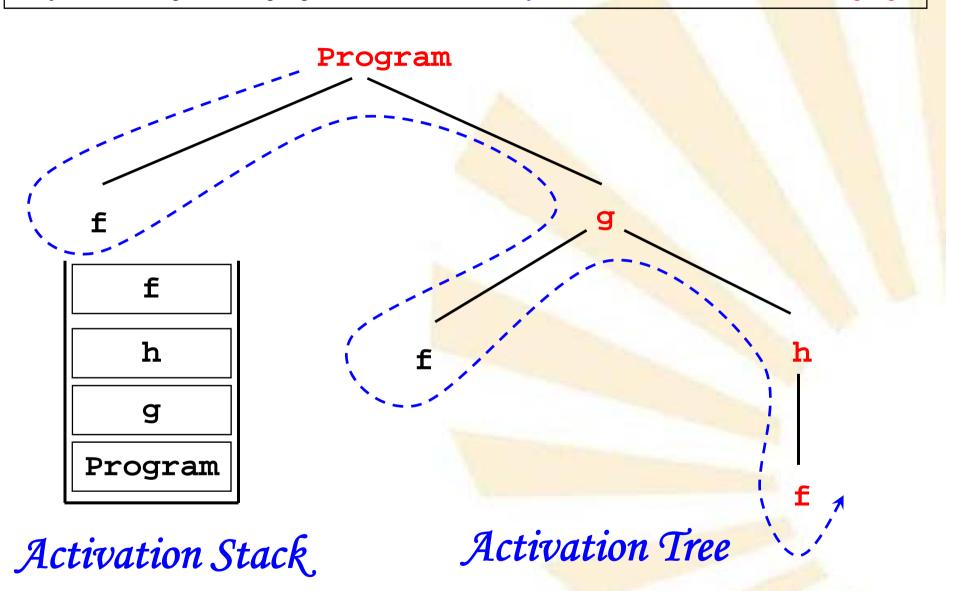
Activation Stack



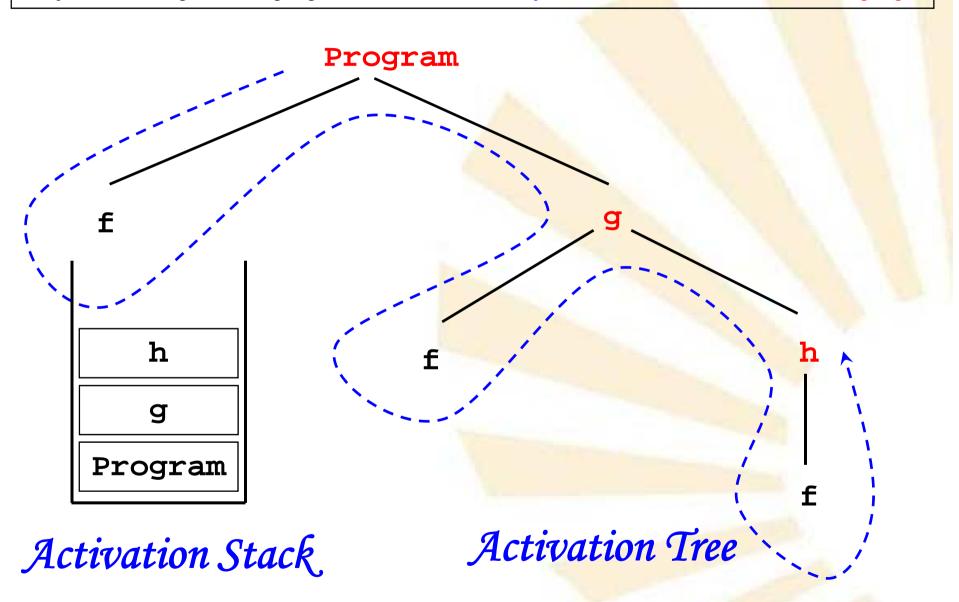
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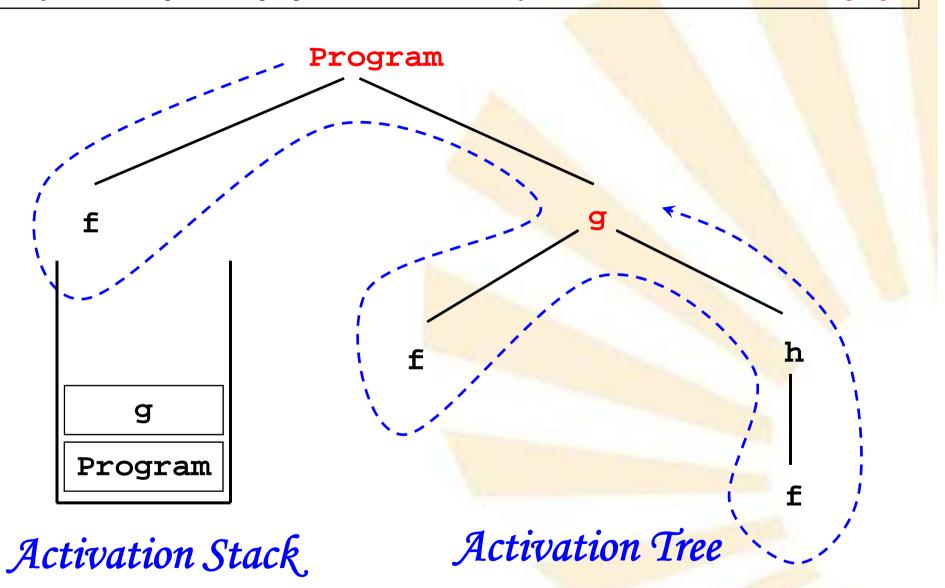


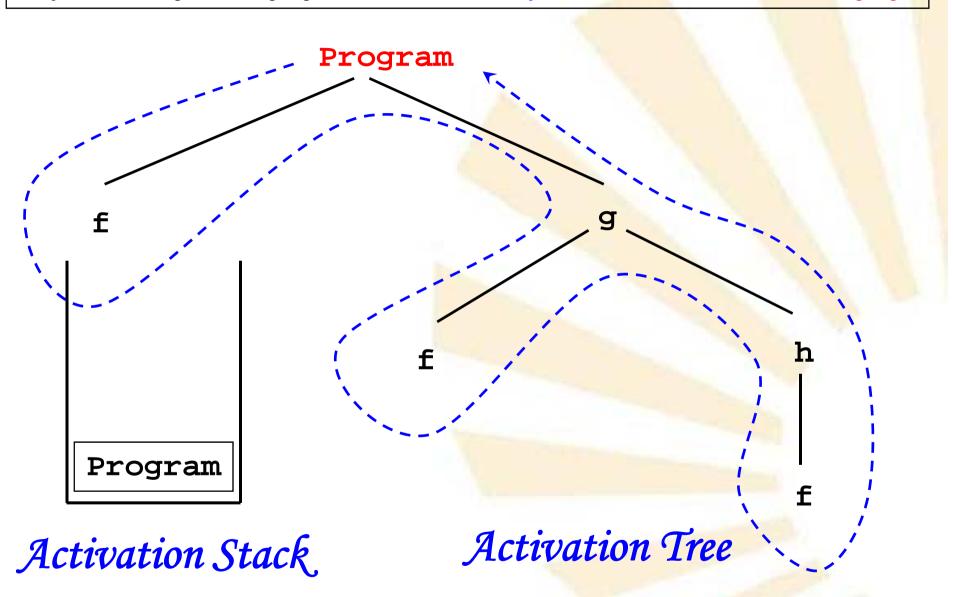
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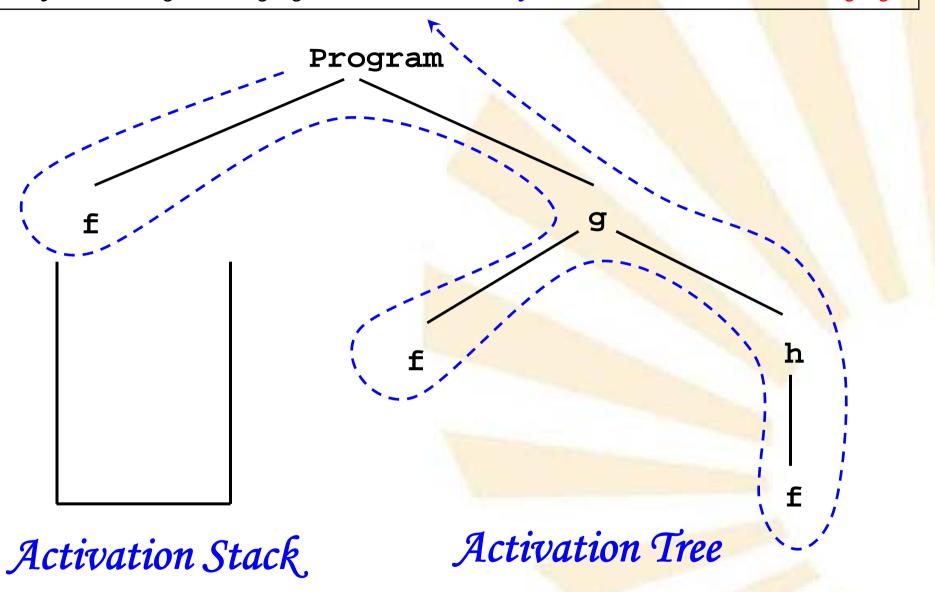


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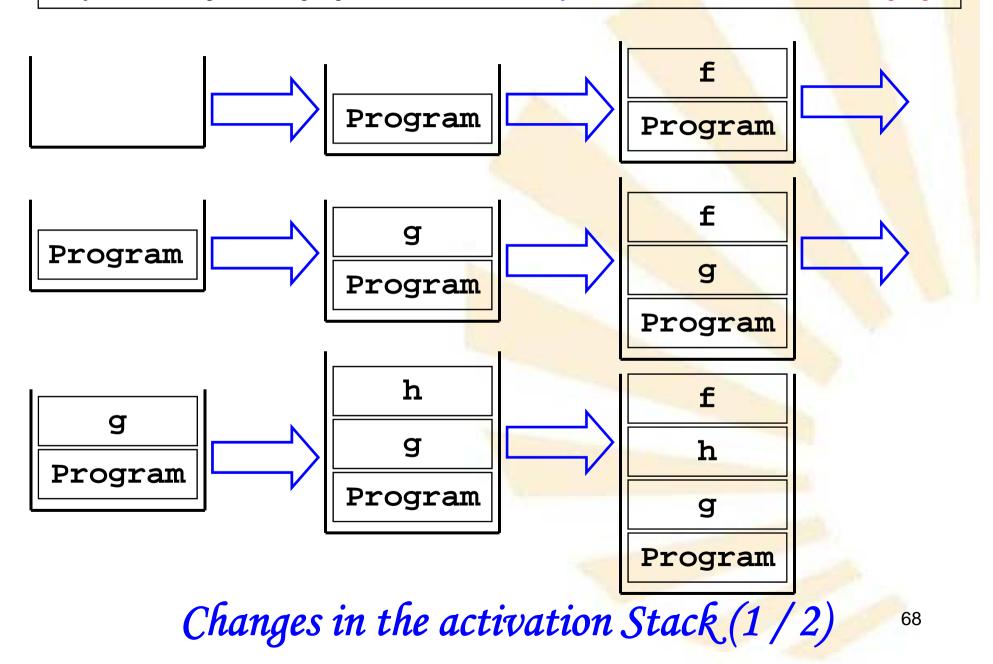




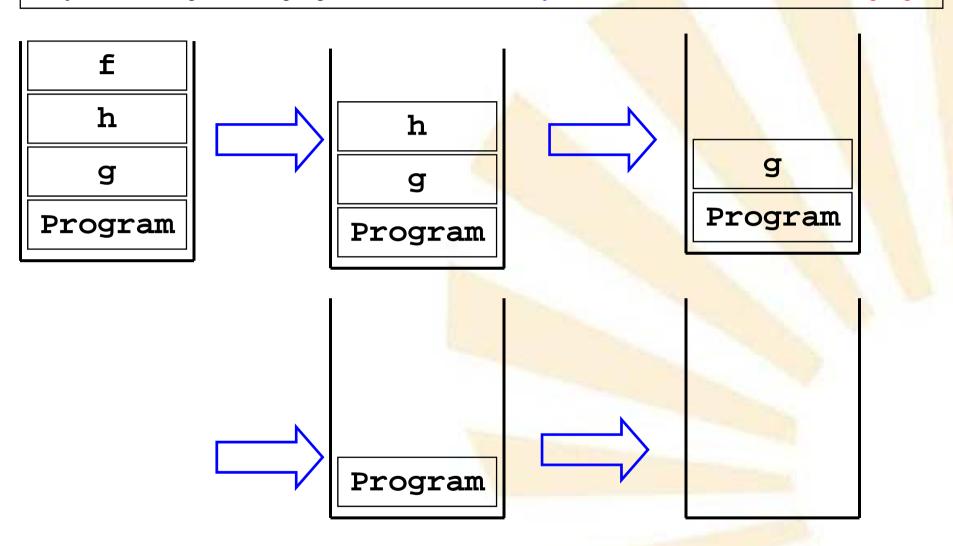




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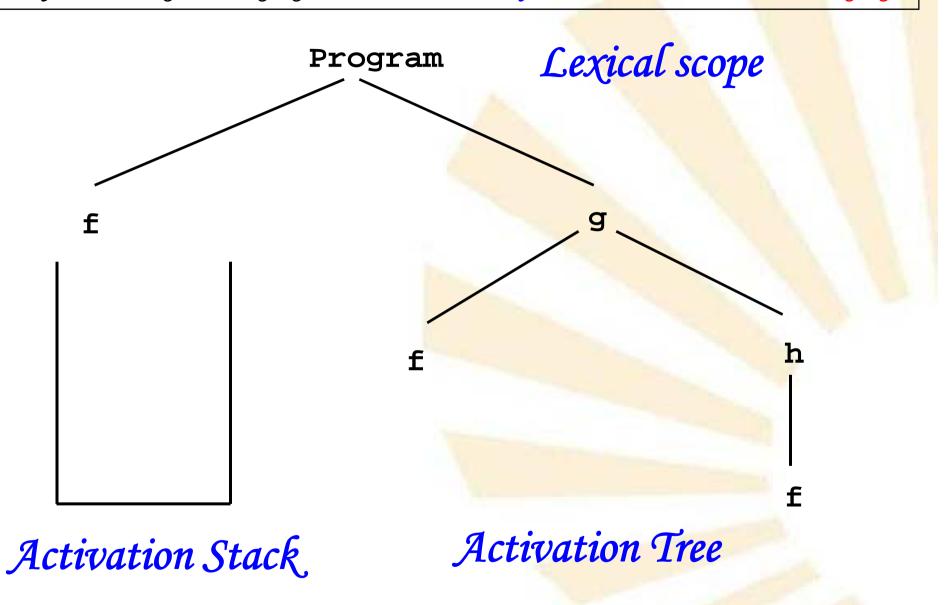
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Changes in the activation Stack (2 / 2)

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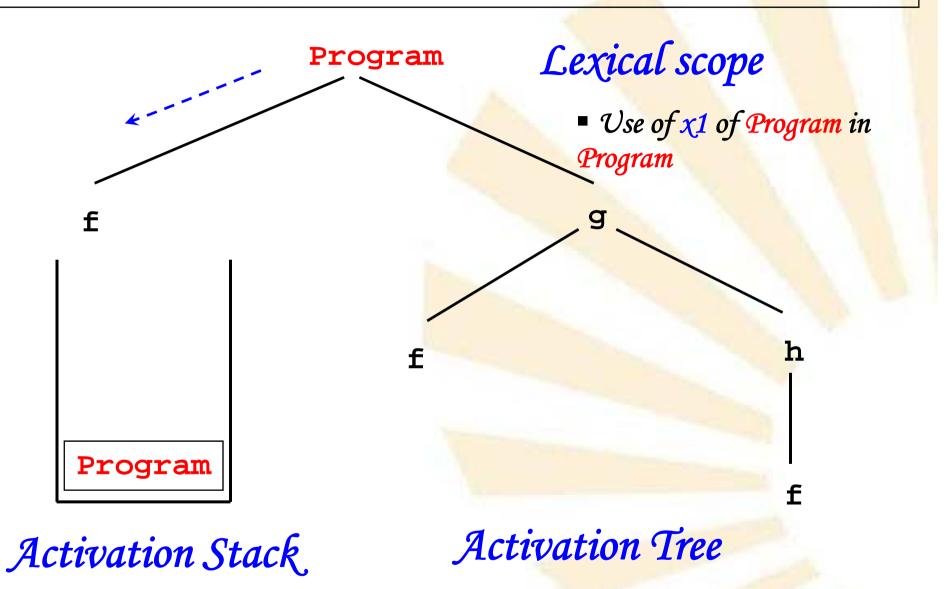
Program **Declaration** of variable **x** (**x**₁) Declaration of procedure f Use of x Declaration of procedure g Run with lexical scope **Declaration** of variable **x** (\mathbf{x}_2) Declaration of procedure h Use of x Call to f Call to f Call to h if condition = true then Call to q else Use of x Use of x Call to f Call to g 70



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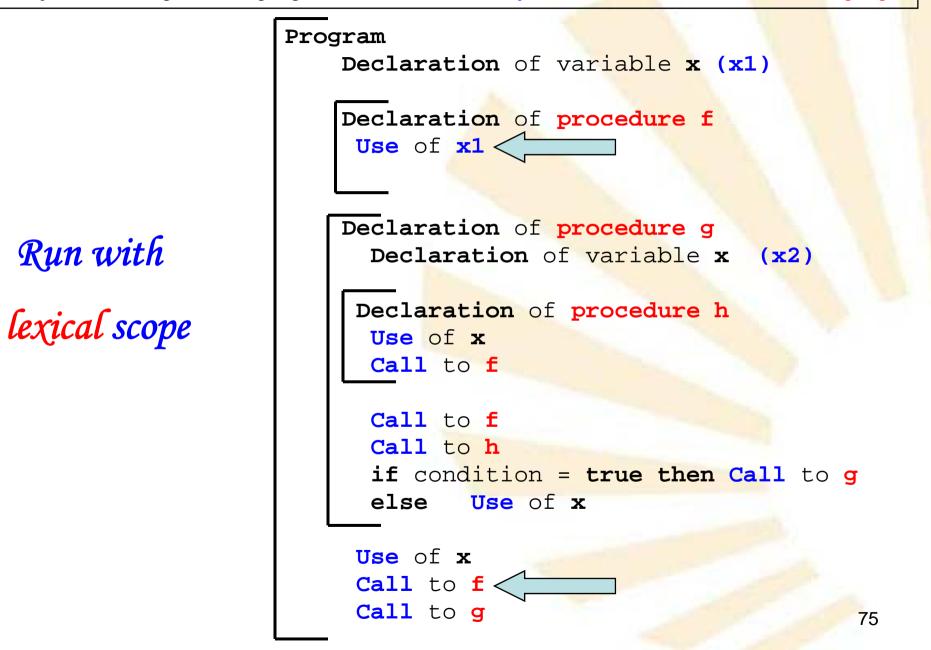
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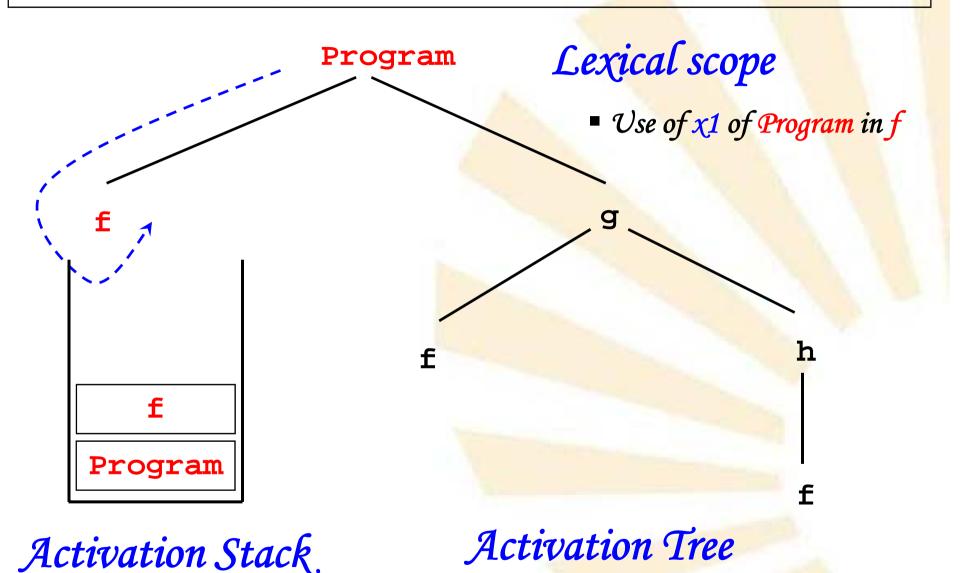


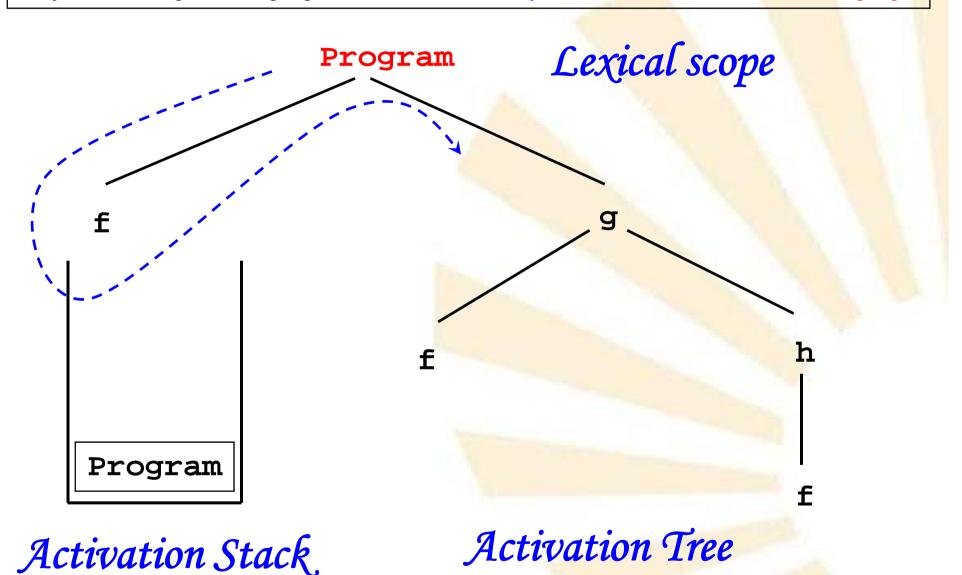
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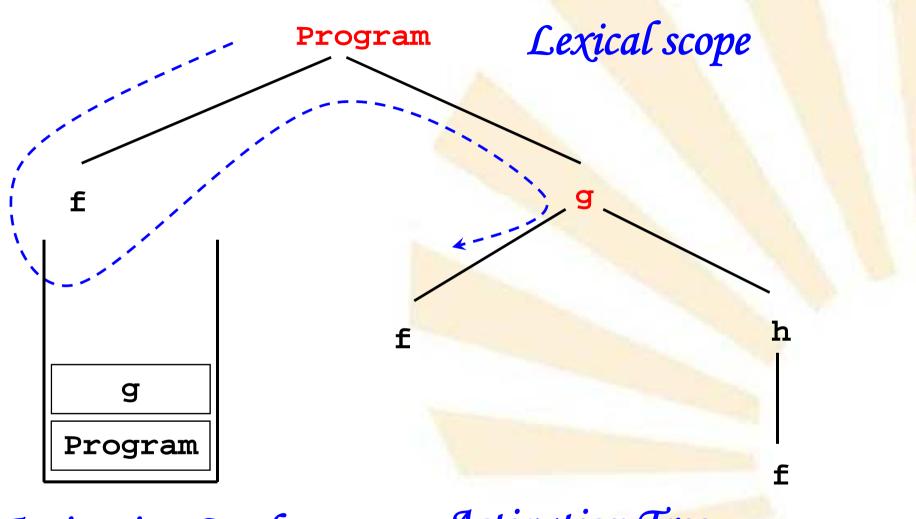






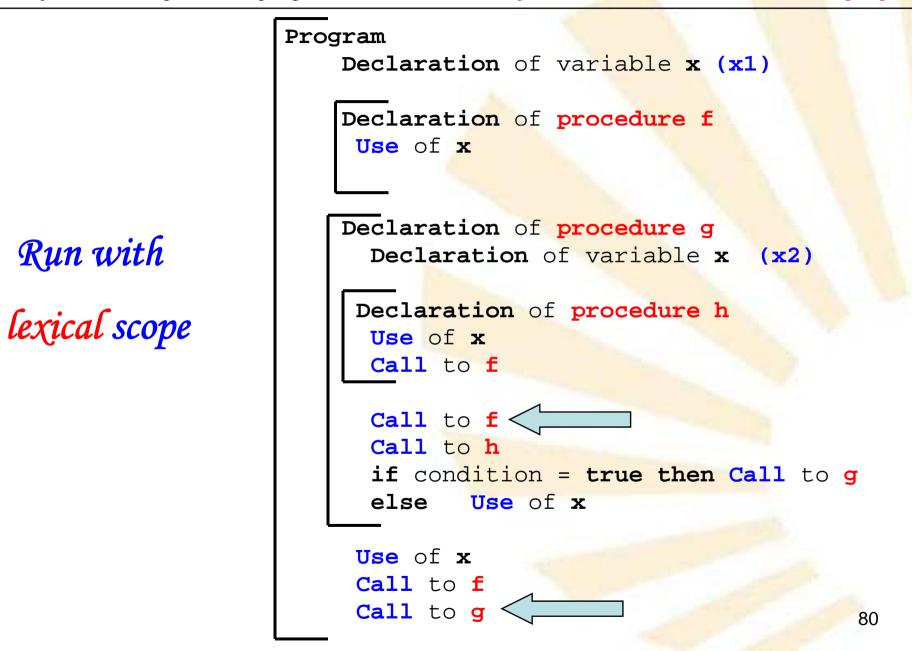
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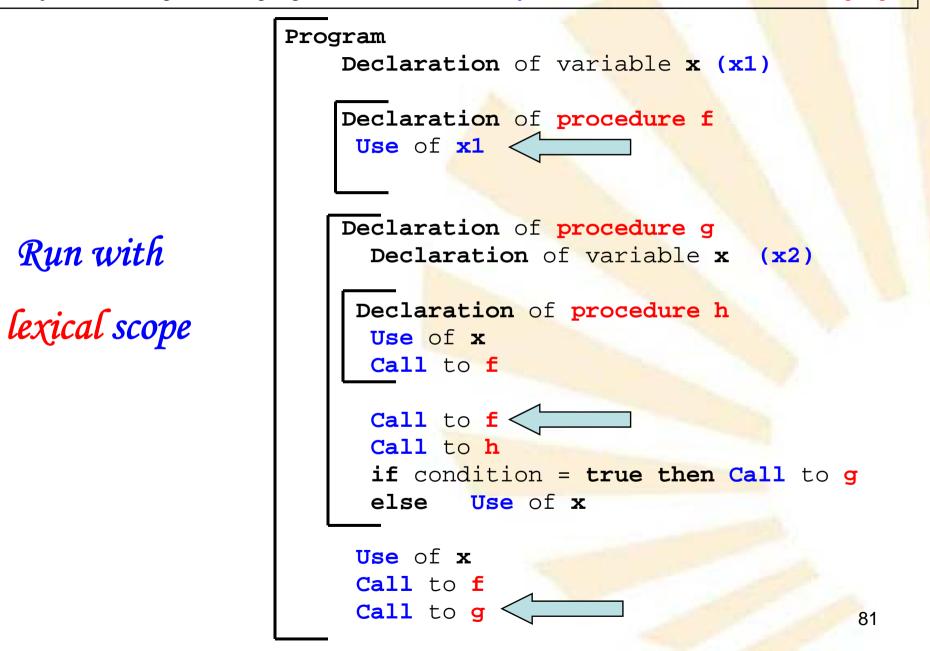
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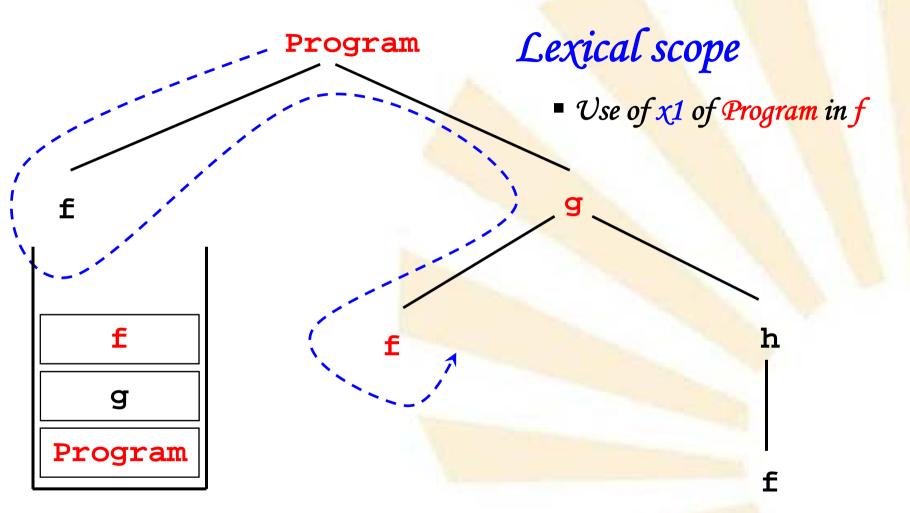


Activation Stack

Activation Tree

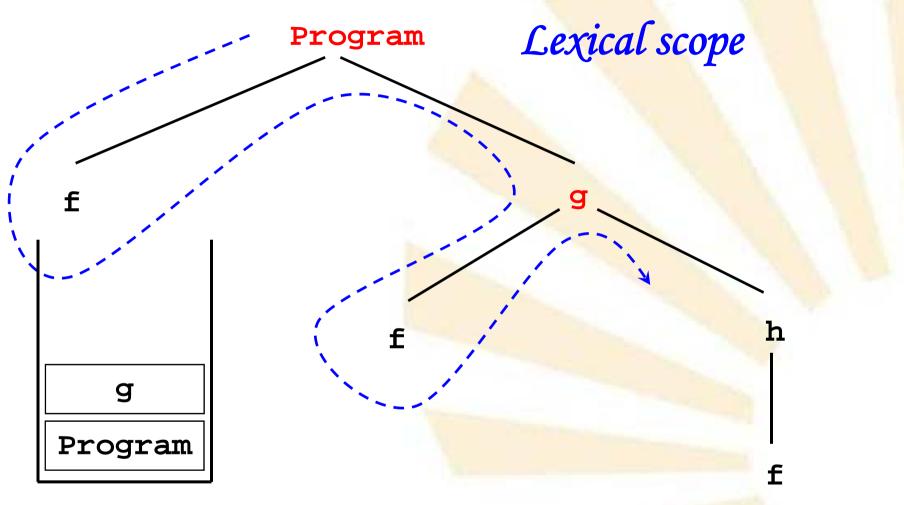






Activation Stack

Activation Tree

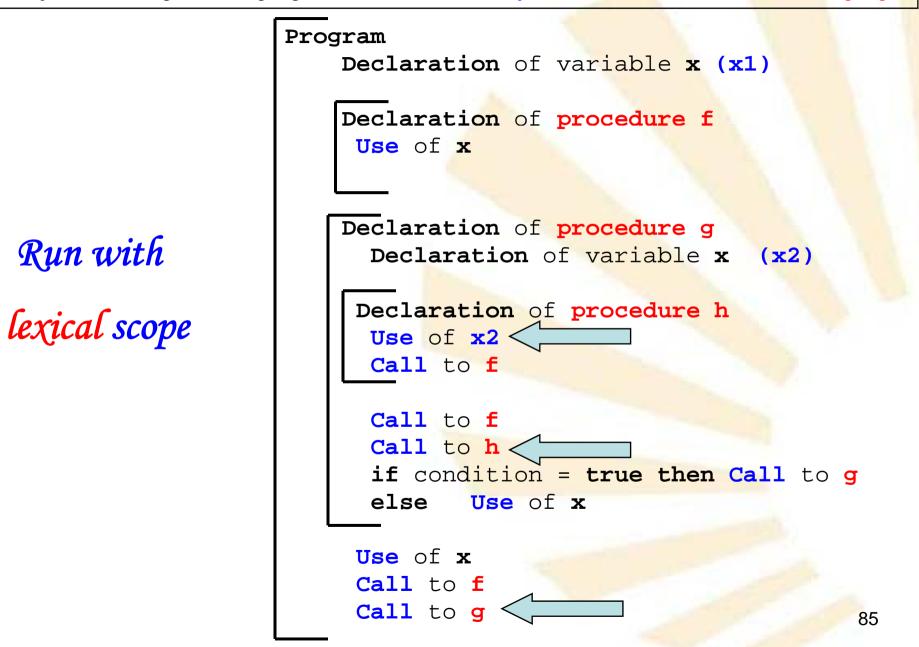


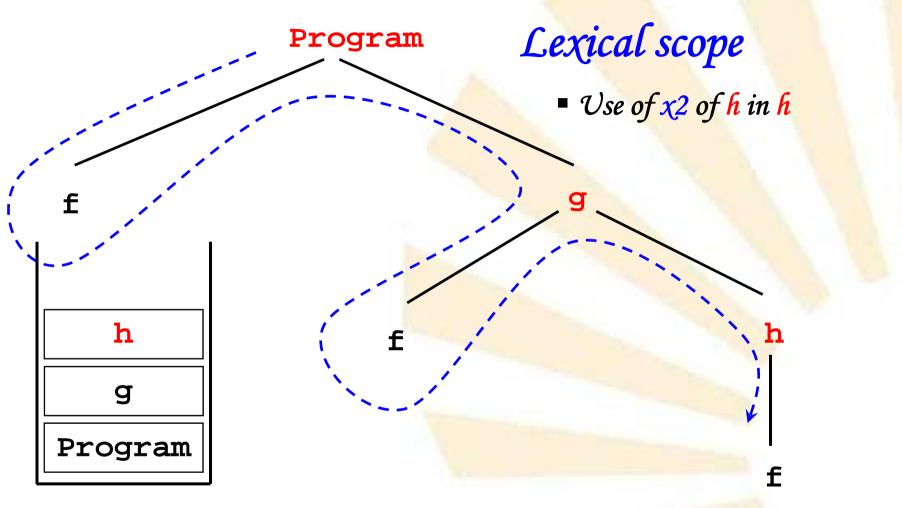
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Activation Tree

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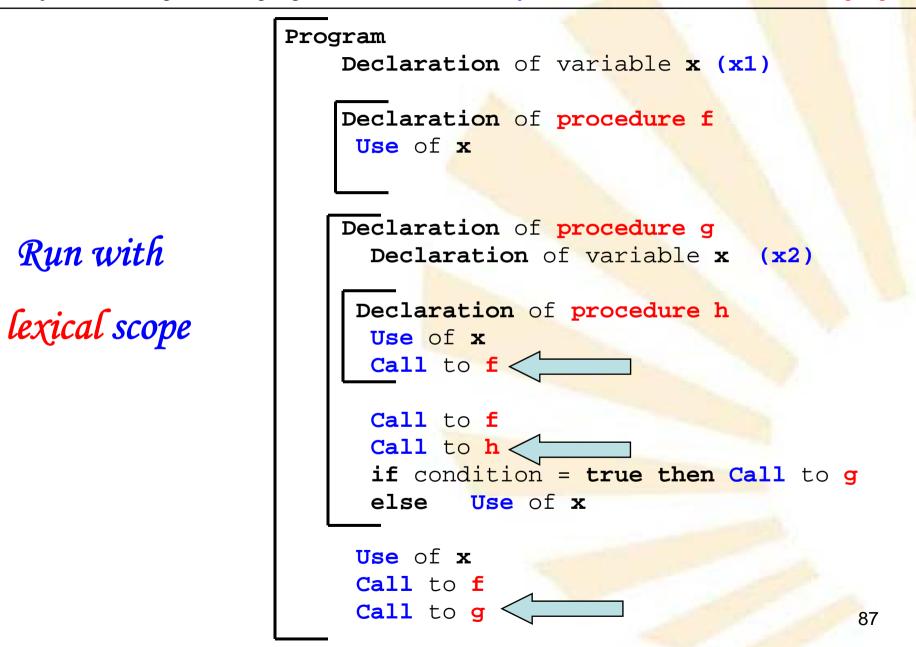
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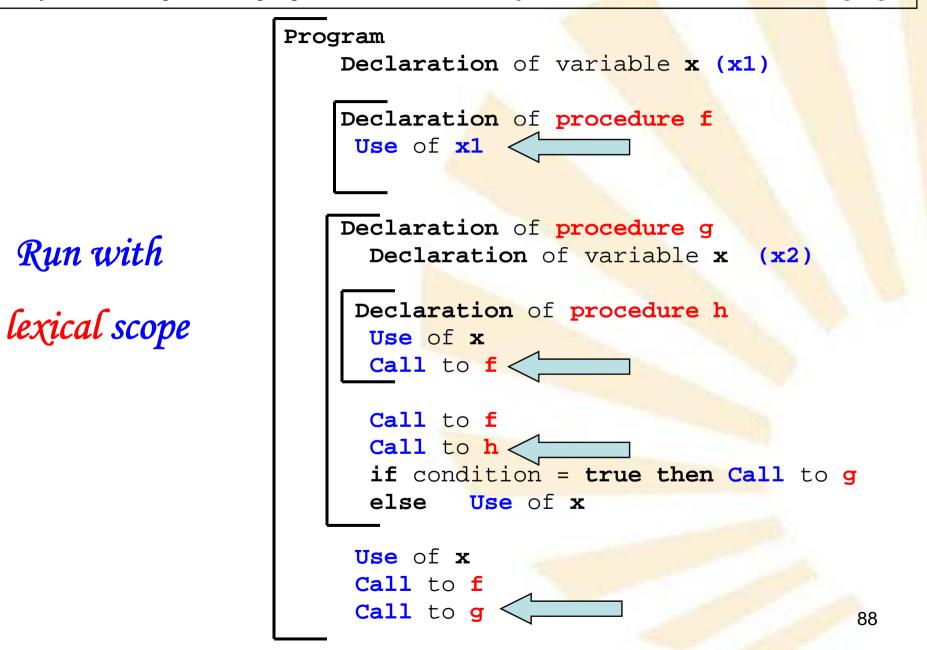


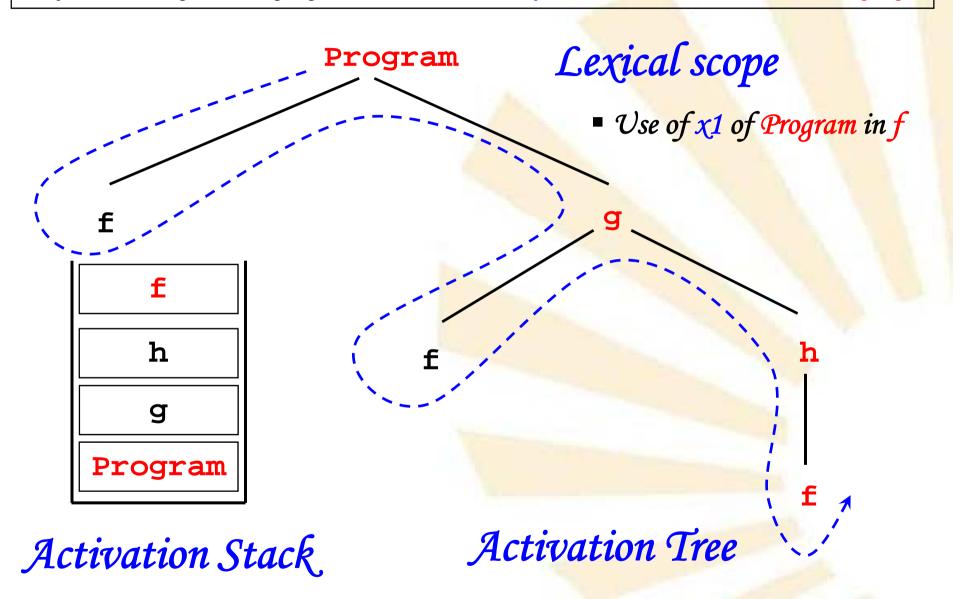


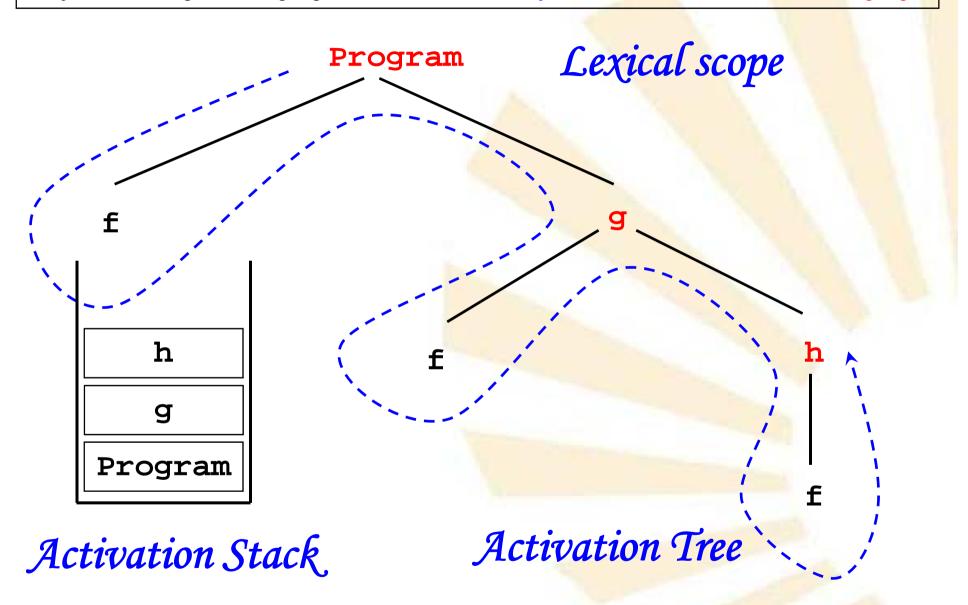
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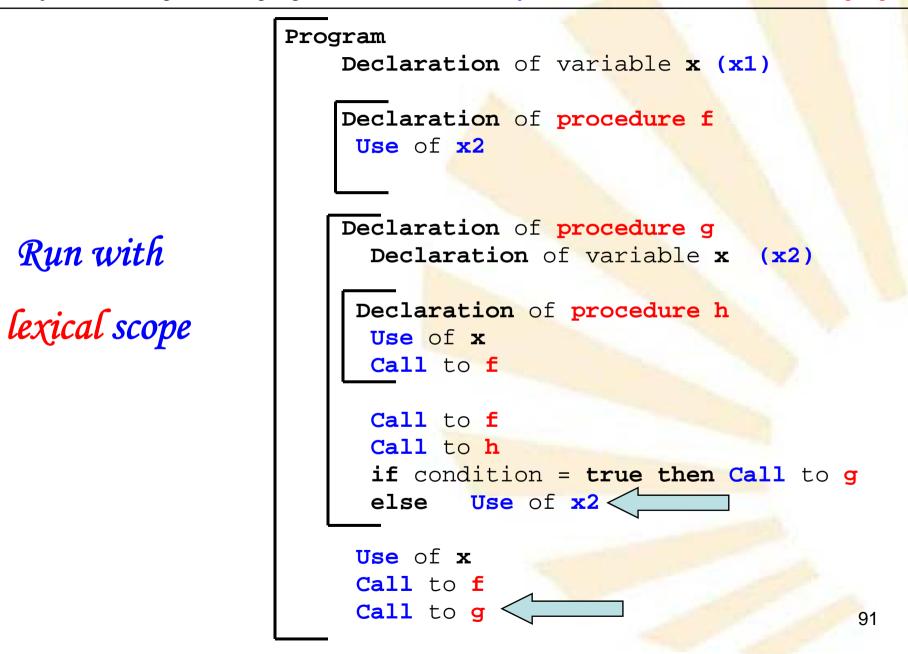
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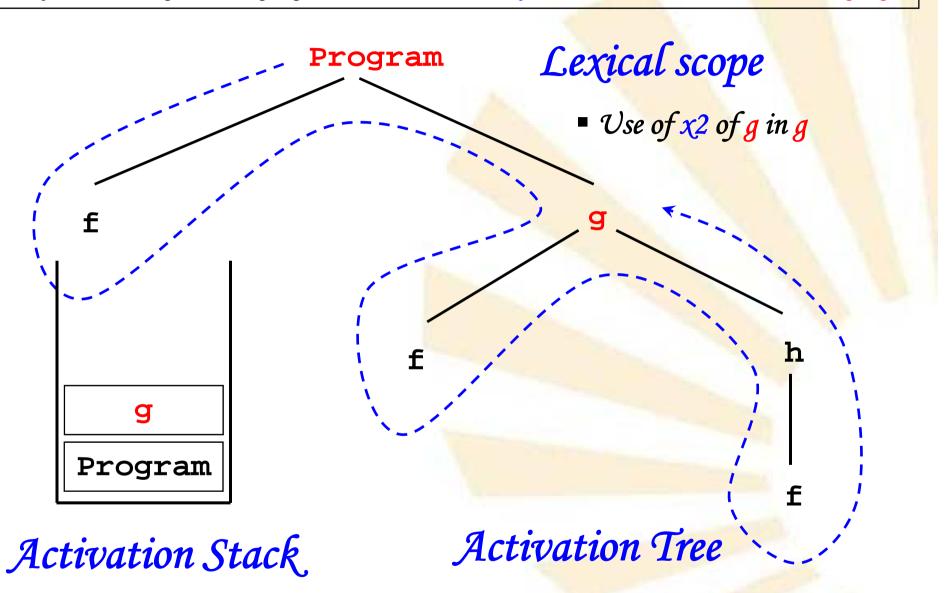






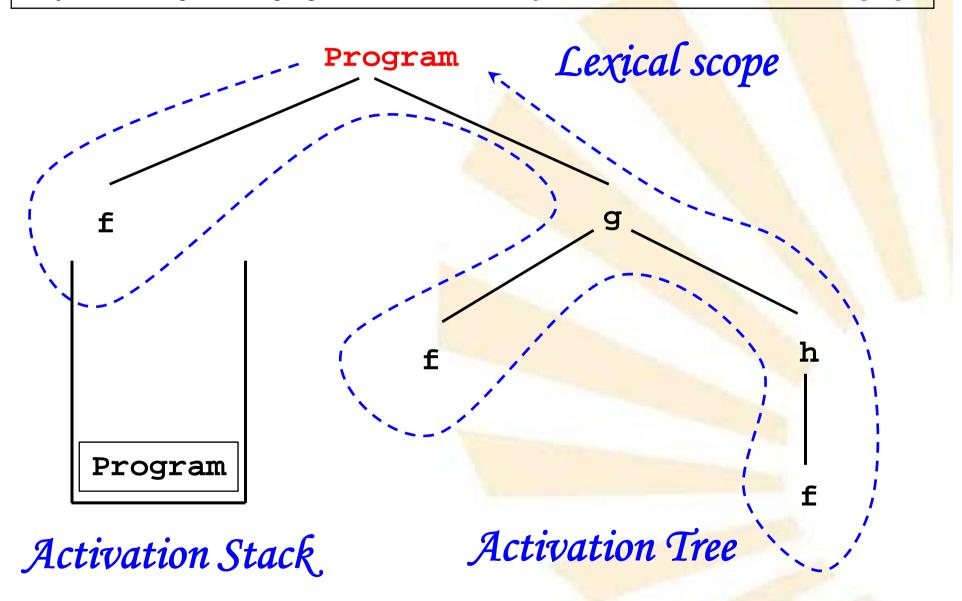


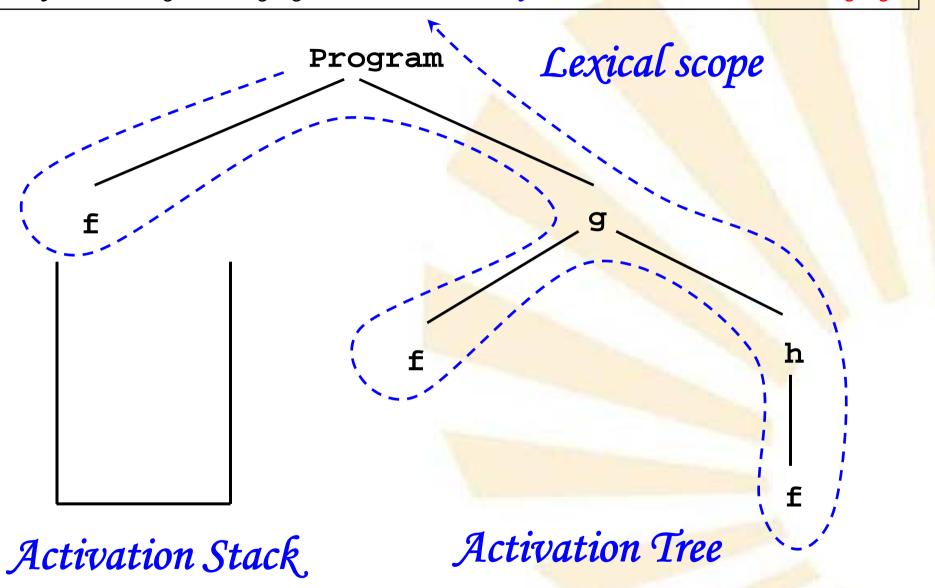




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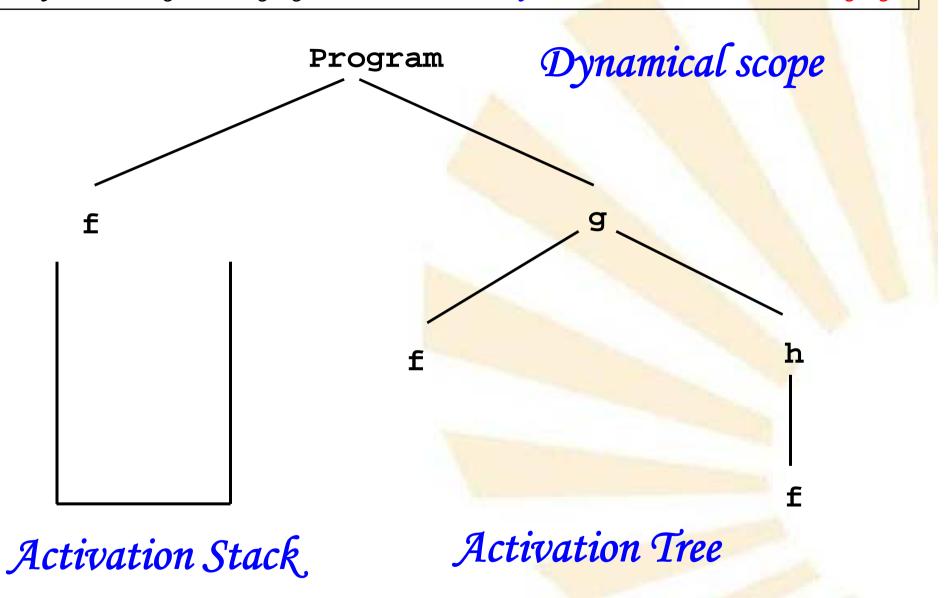
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Subject 1.- Introduction to Scheme language

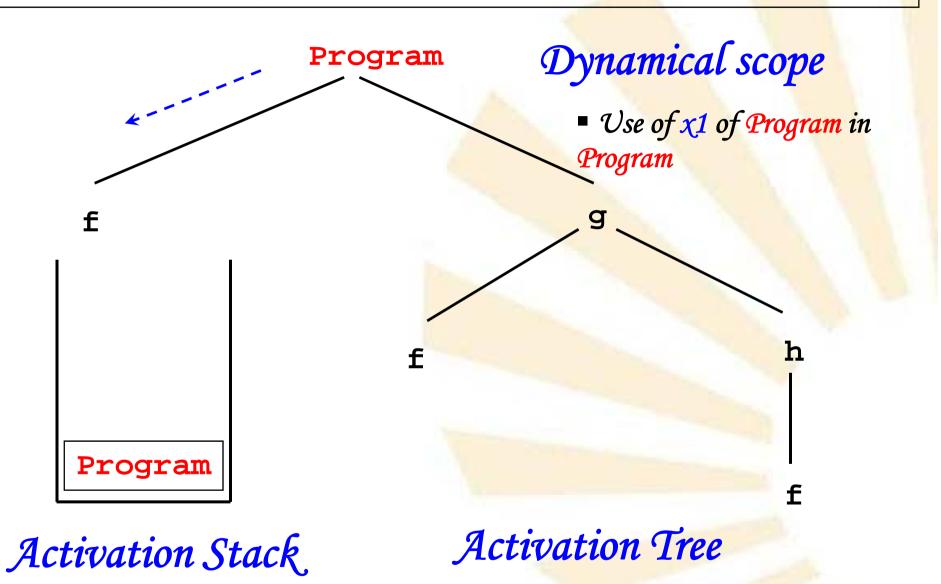
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Subject 1.- Introduction to Scheme language

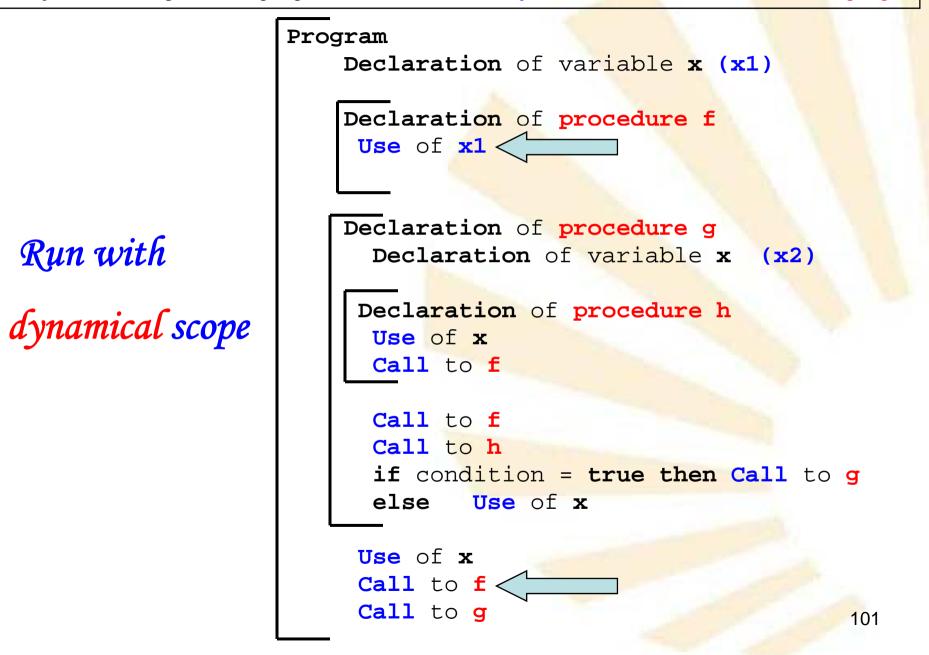
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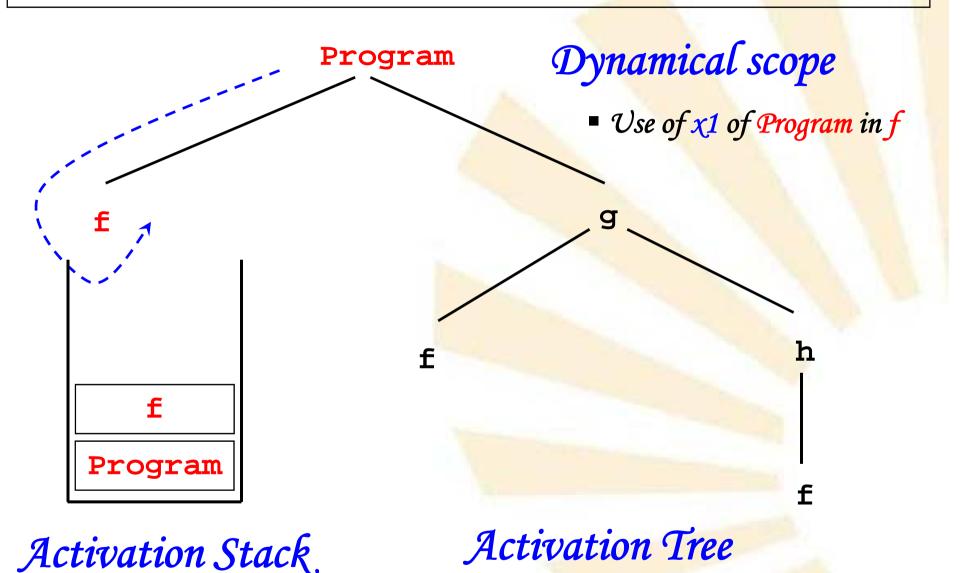


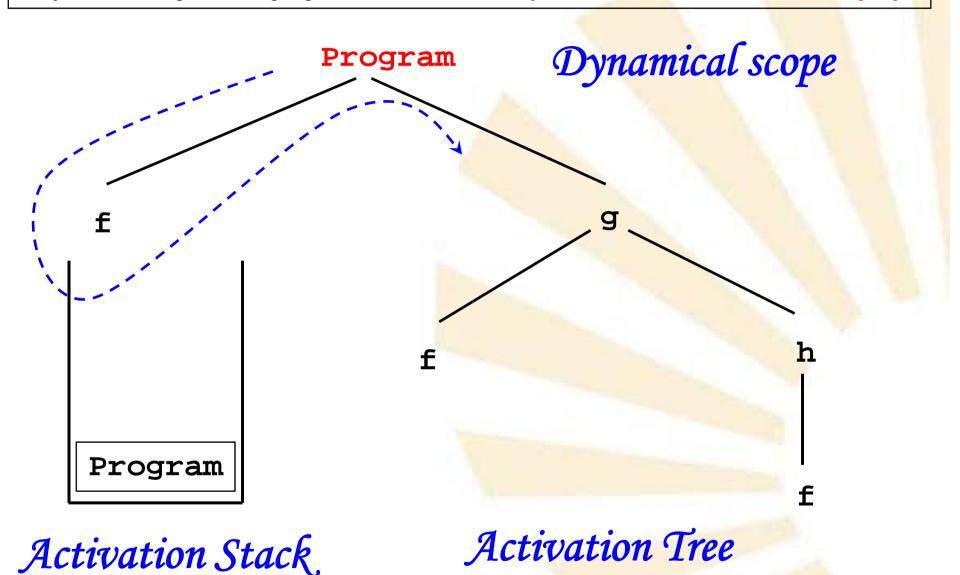
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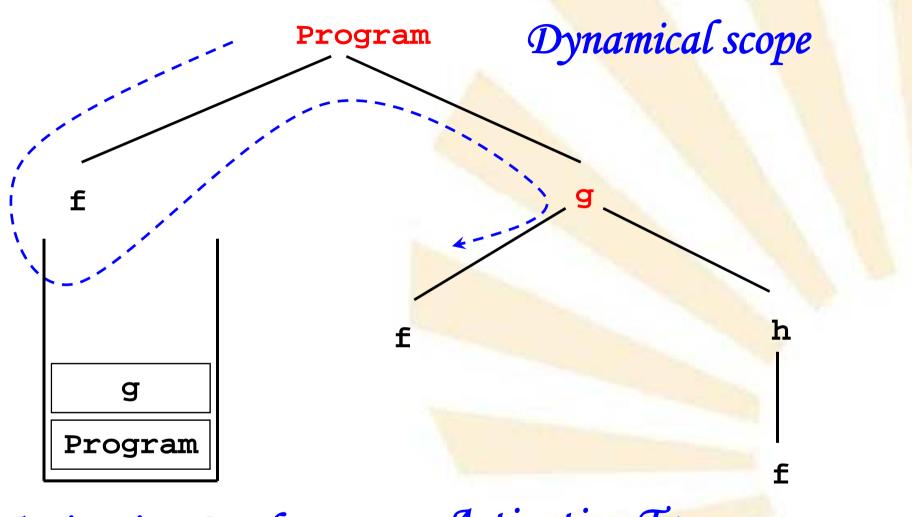




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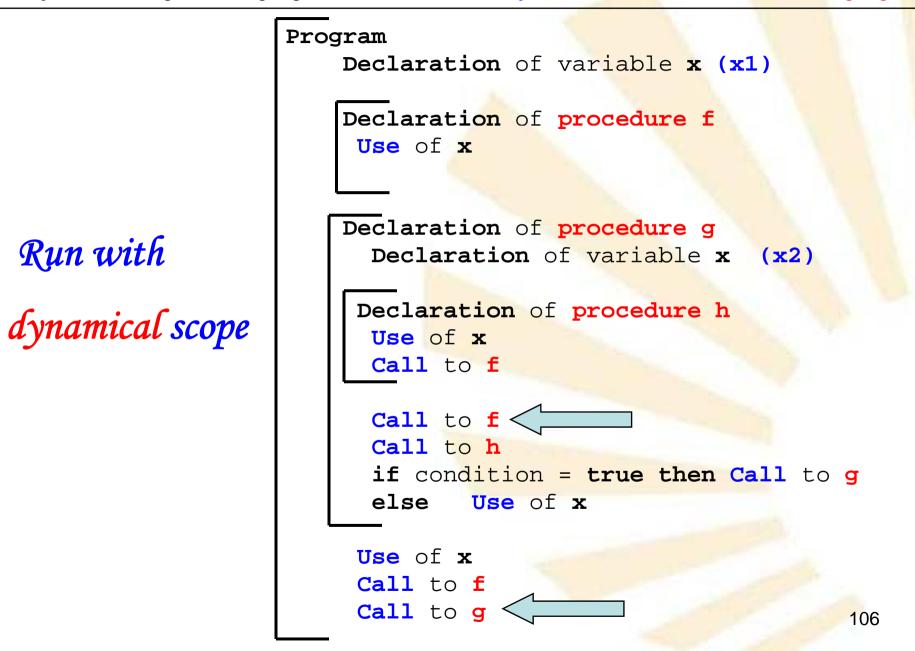
Subject 1.- Introduction to Scheme language

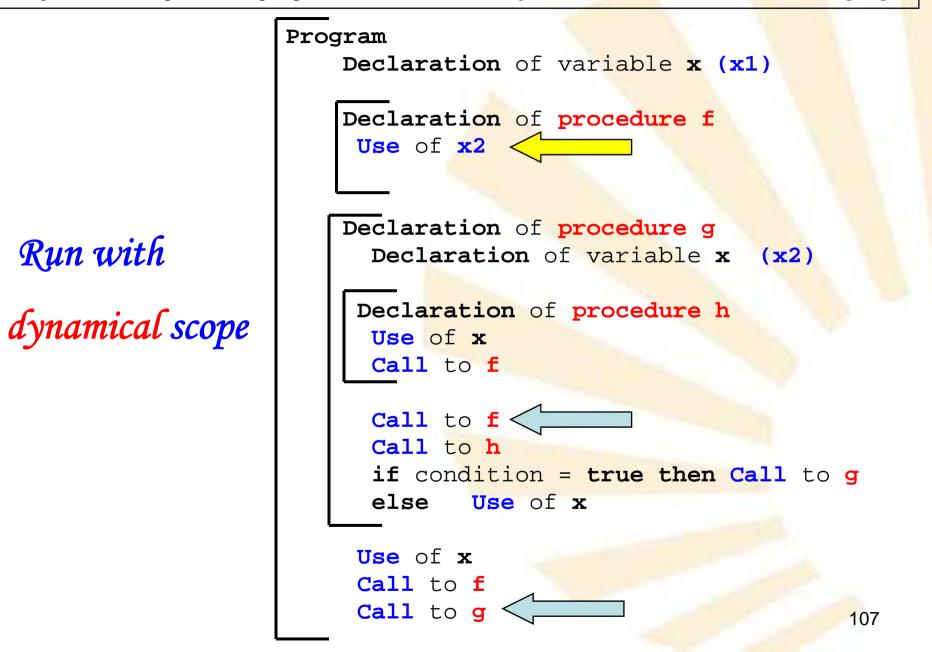
Program **Declaration** of variable **x** (x1) Declaration of procedure f Use of x Declaration of procedure g Run with dynamical scope **Declaration** of variable x (x2) Declaration of procedure h Use of x Call to f Call to f Call to h if condition = true then Call to q else Use of x Use of x Call to f Call to g 104

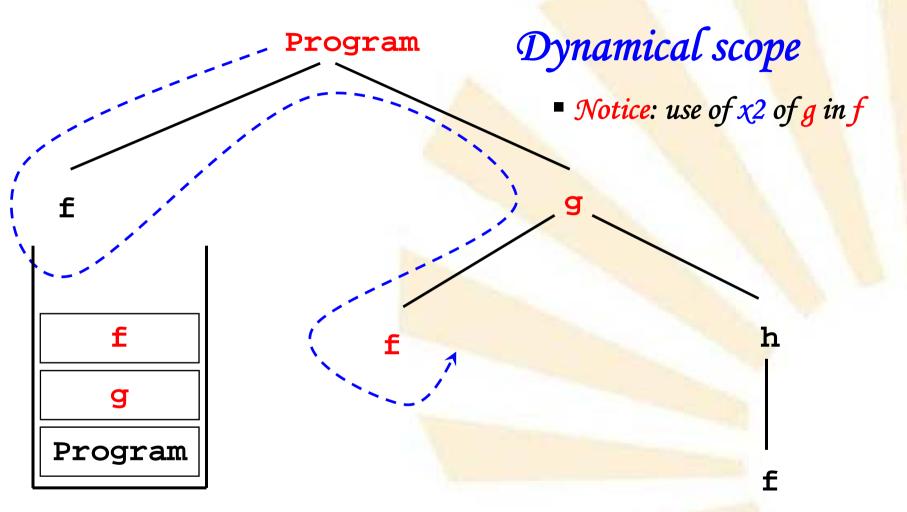


Activation Stack

Activation Tree

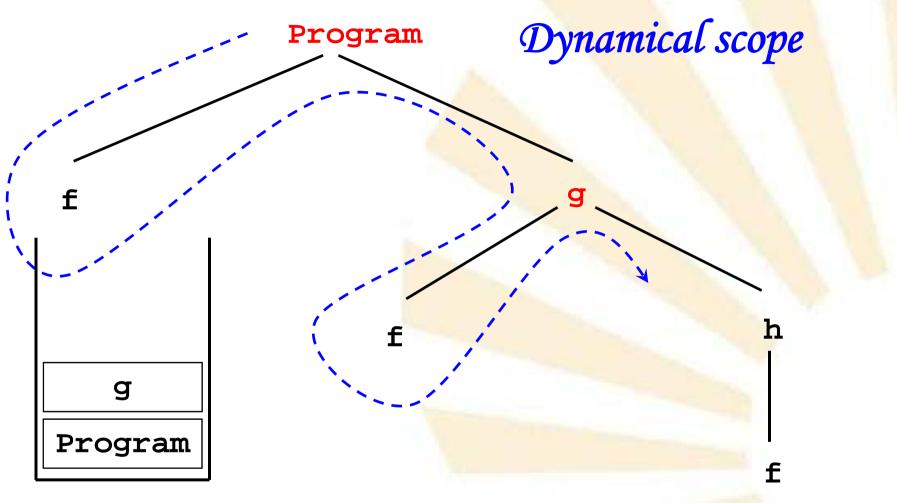






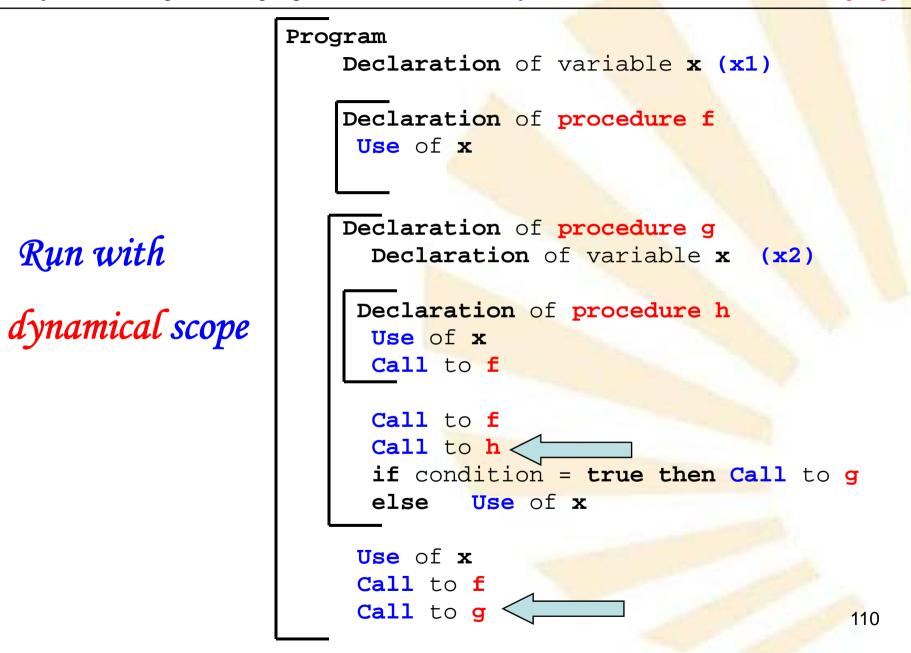
Activation Stack

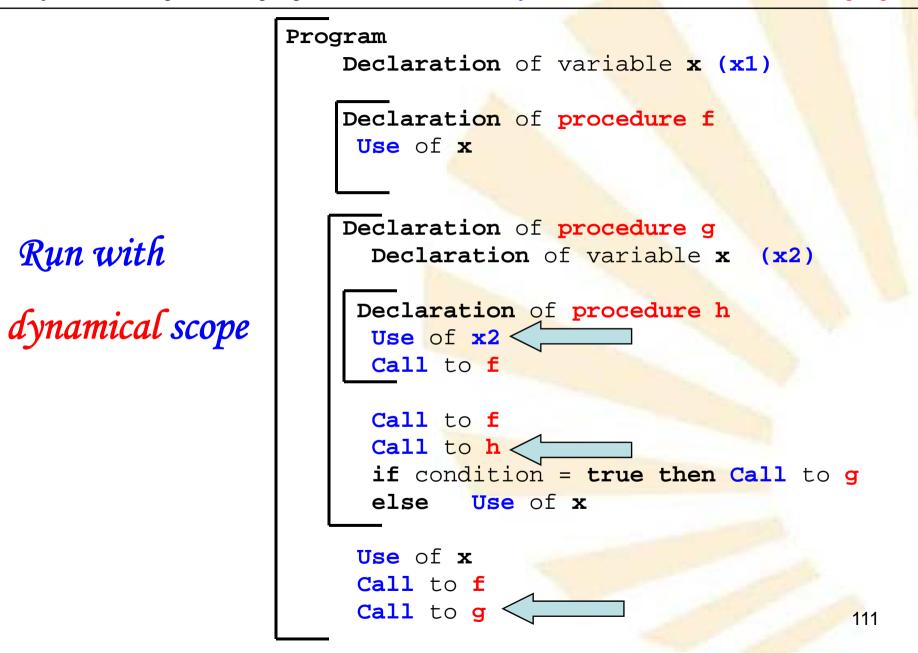
Activation Tree

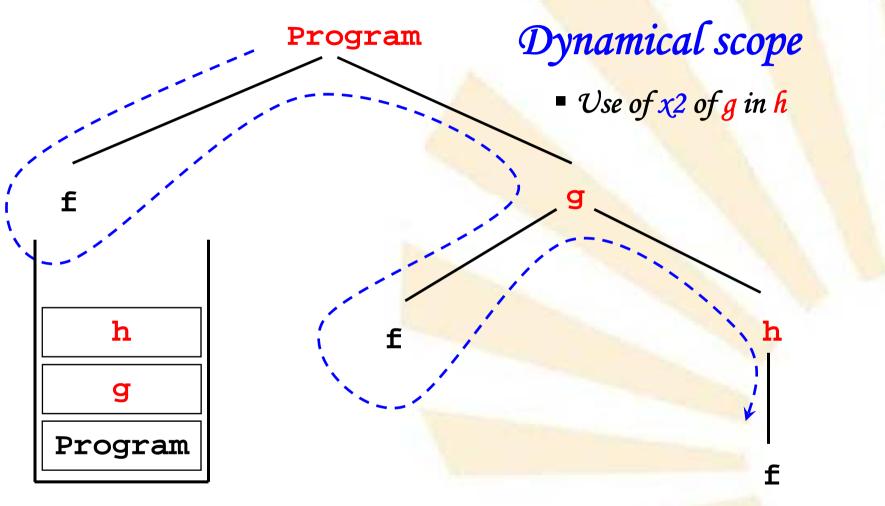


Activation Stack

Activation Tree

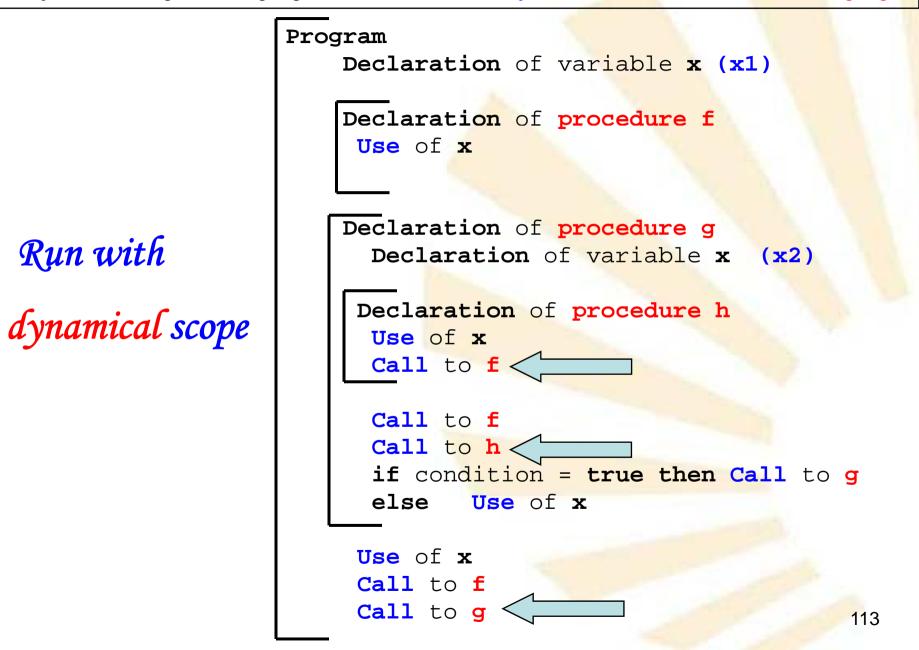


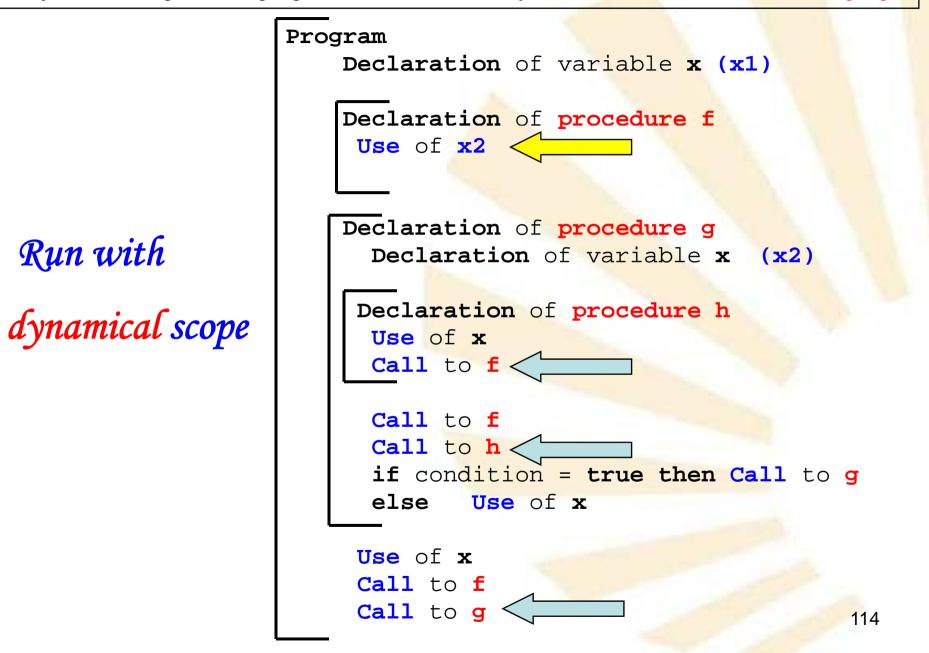


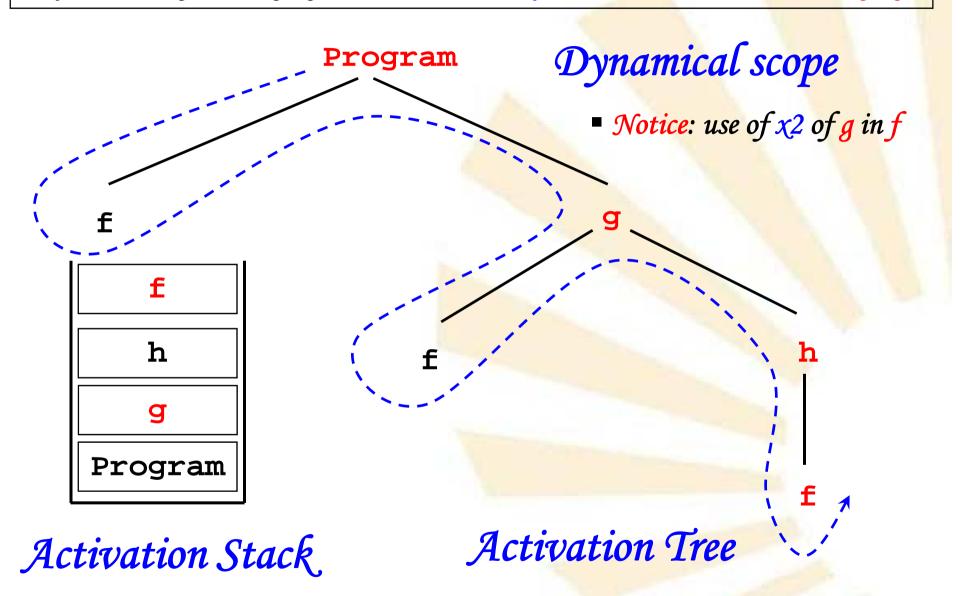


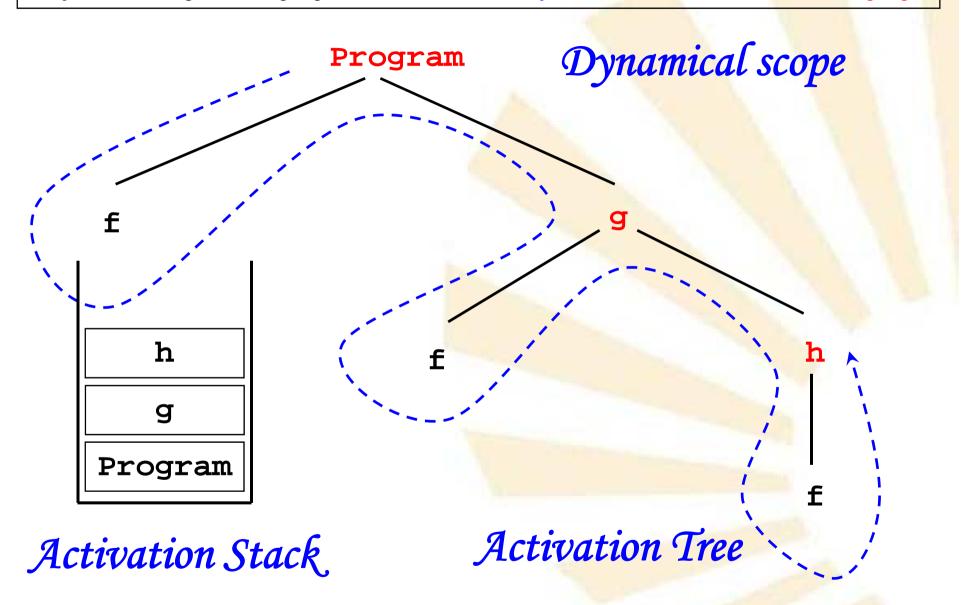
Activation Stack

Activation Tree



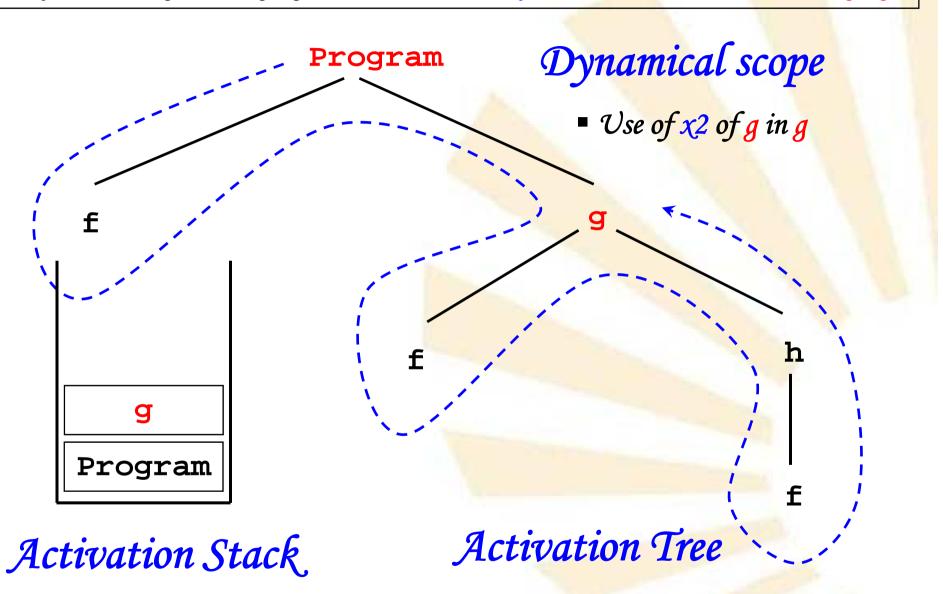






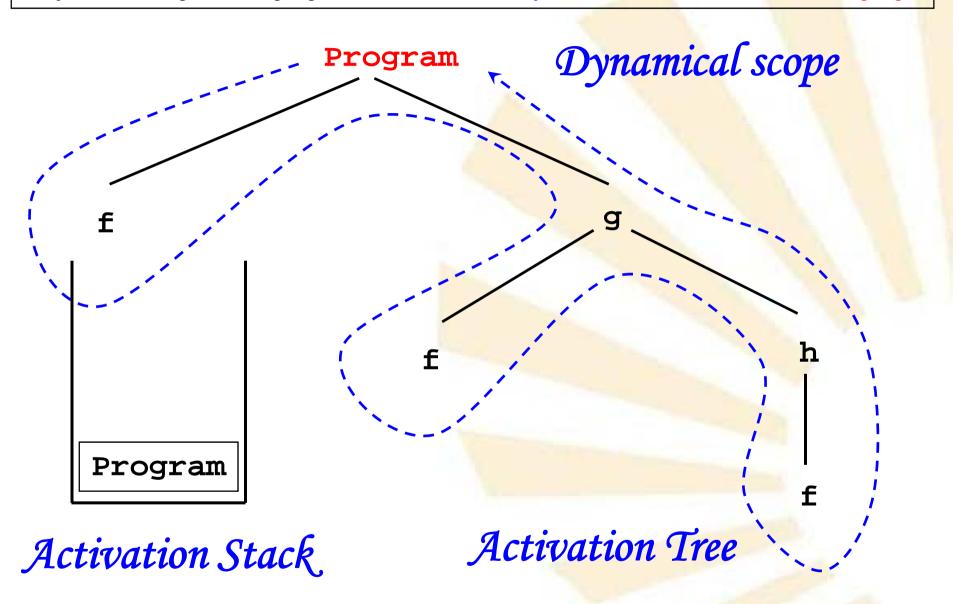
Subject 1.- Introduction to Scheme language

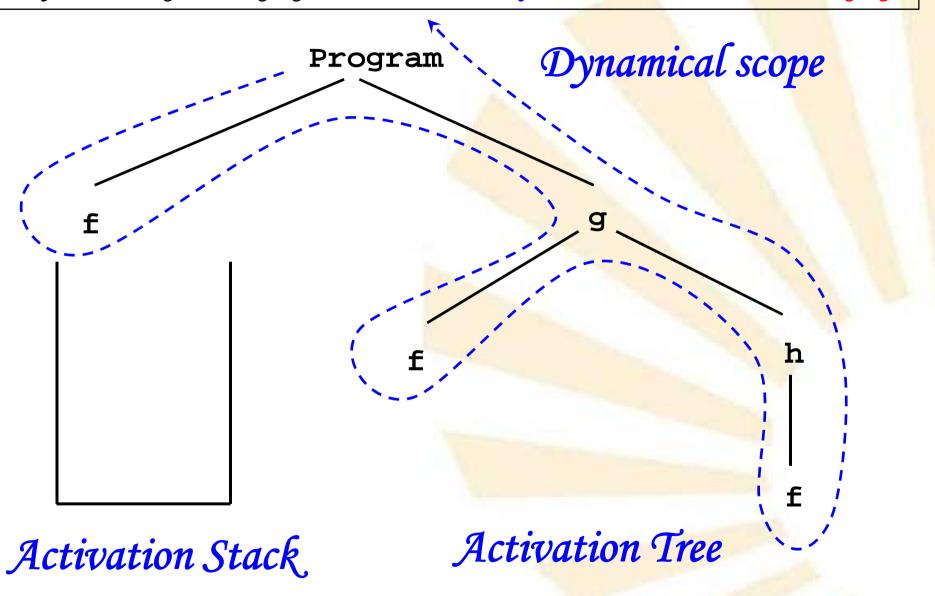
Program **Declaration** of variable **x** (x1) Declaration of procedure f Use of x Declaration of procedure g Run with **Declaration** of variable x (x2) dynamical scope Declaration of procedure h Use of x Call to f Call to f Call to h if condition = true then Call to g else Use of x2 Use of x Call to f Call to g 117



Subject 1.- Introduction to Scheme language

Program **Declaration** of variable **x** (x1) Declaration of procedure f Use of x Declaration of procedure g Run with dynamical scope **Declaration** of variable x (x2) Declaration of procedure h Use of x Call to f Call to f Call to h if condition = true then Call to q else Use of x2 Use of x Call to f Call to g 119





- 2. Historic Summary of Scheme
 - ✓ LISP
 - ✓ Compilation versus Interpretation
 - ✓ Dynamically versus statically scope
 - ✓ Origin of Scheme

- 2. Historic Summary of Scheme
 - ✓ Origin of Scheme:
 - Gerald Jay Sussman (MIT) and Guy Lewis Steele Jr.
 - Question:

How would LISP be with lexical or static scope rules?

- ➤ Answer: new language → Scheme
 - More efficient implementation of recursion
 - First class functions.
 - Rigorous semantic rules
- Influence on Common LISP: lexical scope rules
- Revised ⁵ Report on the Algorithmic Language Scheme

- 2. Historic Summary of Scheme
 - ✓ *Scheme*:
 - Structure of scheme programs
 - Sequence of
 - definitions of functions and variables
 - and *expressions*



CÓRDOBA U<mark>NIVERSIT</mark>Y

SUPERIOR POLYTECHNIC SCHOOL

DEPARTAMENT OF COMPUTER SCIE<mark>NCE AND NUMERICAL AN</mark>ALYSIS



ARTIFICIAL INTELLIGENCE LANGUAGES

TECHNICAL ENGINEERING IN MANAGEMENT COMPUTER SCIENCE

TECHNICAL ENGINEERING IN SYSTEMS COMPUTER SCIENCE

SECOND COURSE

FIRST FOUR-MONTH PERIOD

ACADEMIC YEAR: 2009 - 2010

